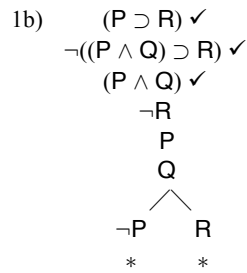
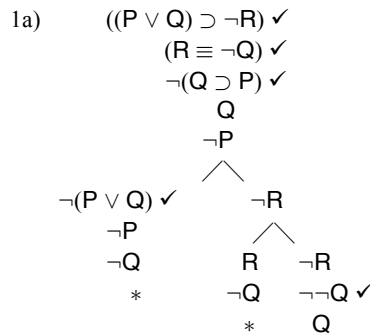


Assignment #5-key



The completed tree closes.
 So the argument is *tautologically valid*.

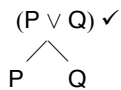


The completed tree has an open branch.
 So the argument is *tautologically invalid*.

2a) If there is a complete tree for the PLC wff A all of whose branches are open, then A is a tautology.

Comment:

This is *not* a true claim about PLC trees. A branch on a complete tree represents an attempt to make the wff(s) on its trunk true. A closed branch represents a failed attempt, while an open branch represents a successful attempt. In other words, an open branch on a complete tree represents a *valuation* that makes the wff(s) on the trunk true. But the number of branches on a completed tree are not necessarily equal to the total number of possible valuations of the wff(s) on the trunk. There can be more valuations of the wff(s) on the trunk than there are branches in the completed tree. So just because all branches on a complete tree with a single wff on its trunk are open doesn't necessarily mean all valuations of the wff make it true. There could be one or more valuations that make it false. Here is an example:



All branches of this completed tree are open, but of course the wff $(P \vee Q)$ is not a tautology!

2b) If there is a complete tree for the PLC wff A all of whose branches close, then A is a contradiction.

Proof:

The tree rules *guarantee* that if a valuation that makes the wff(s) on the trunk true exists, then it will be found in the form of an open branch (the tree rules require you to unpack a wff into its "truth-makers"). So if there are no open branches on a complete tree with a single wff on its trunk, then there are no valuations that make that wff true. So the wff must be a contradiction.

2c) If there is a complete truth tree for the PLC wff $(A \supset B)$ all of whose branches are closed, then A is a tautology and B is a contradiction.

Proof:

Suppose: There's a complete truth tree for $(A \supset B)$ all of whose branches close.

Then: $(A \supset B)$ is a contradiction (see #2b).

So: All valuations make $(A \supset B)$ false.

So: All valuations make A true and B false.

So: A is a tautology and B is a contradiction.

3. For any wffs A_1, \dots, A_n, C , if $\models ((A_1 \wedge \dots \wedge A_n) \supset C)$, then $A_1, \dots, A_n \models C$.

Proof:

Suppose: $\models ((A_1 \wedge \dots \wedge A_n) \supset C)$

Then: No valuation makes $((A_1 \wedge \dots \wedge A_n) \supset C)$ false.

So: No valuation makes $(A_1 \wedge \dots \wedge A_n)$ true and C false.

So: No valuation makes A_1, \dots, A_n true and C false.

So: If any valuation makes A_1, \dots, A_n true, it also makes C true.

Thus: $A_1, \dots, A_n \models C$.