

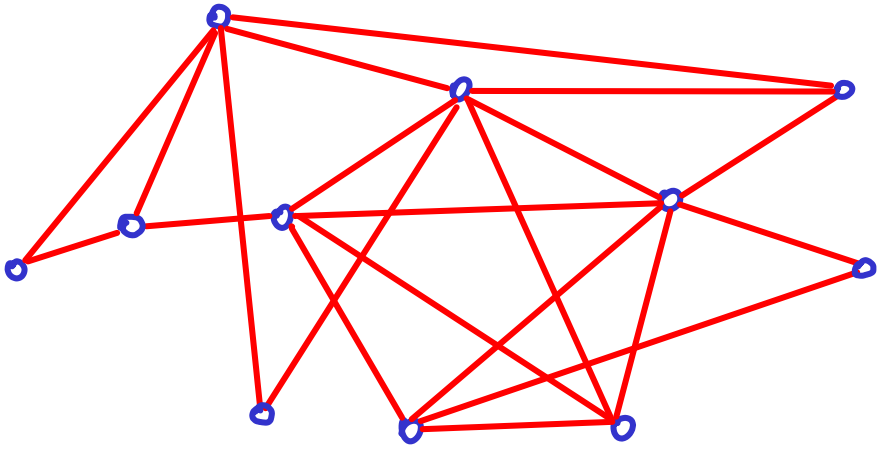
Game for 2 players:

Draw a complete graph,  
each player can draw edges using one color,  
take turns coloring one edge at a time.

Whoever completes a triangle first wins.

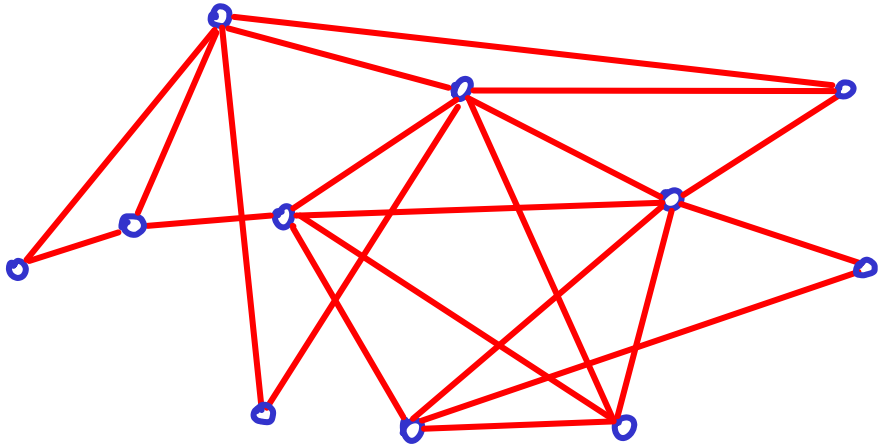
Is there always a winner?

# Cliques



Given  $G$ , a subset  $S$  of  $V(G)$  is a clique if every  $s_i, s_j \in S$  share an edge in  $G$ .

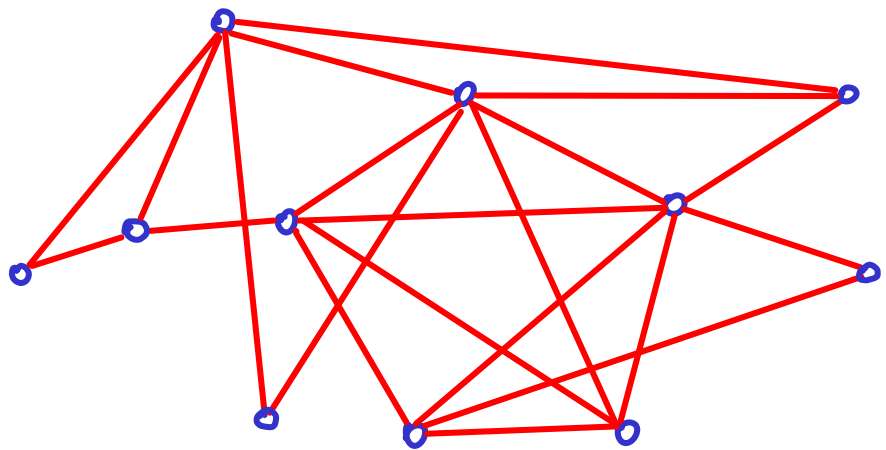
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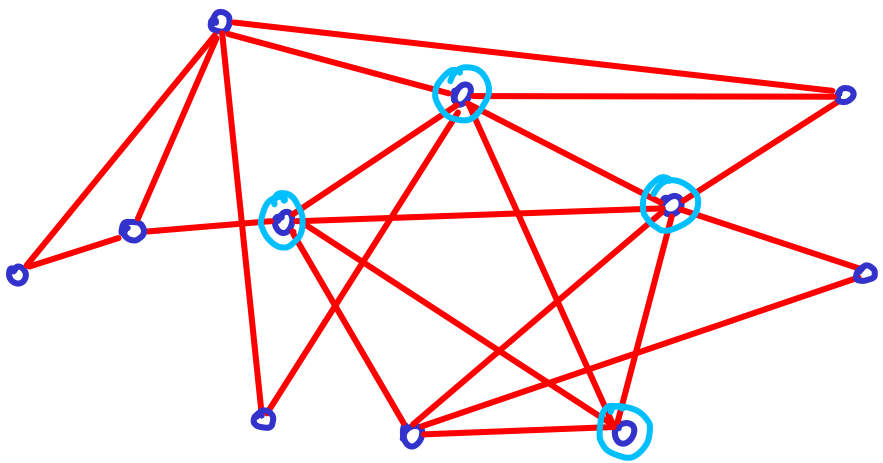


How many cliques?  
What is the largest clique?

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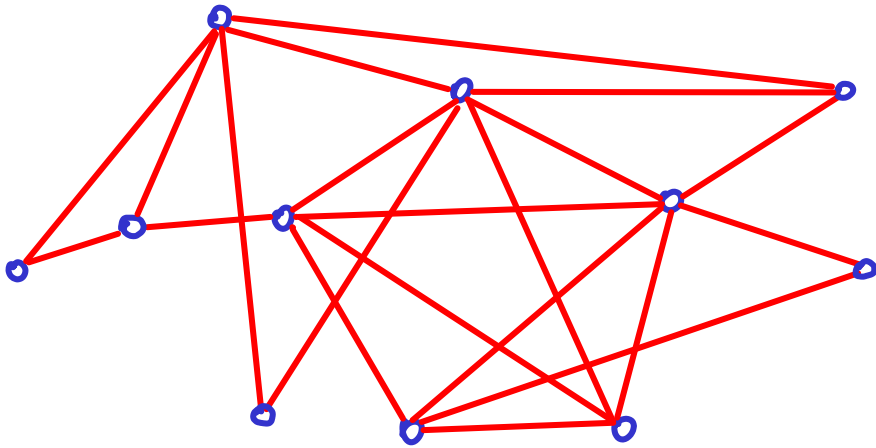


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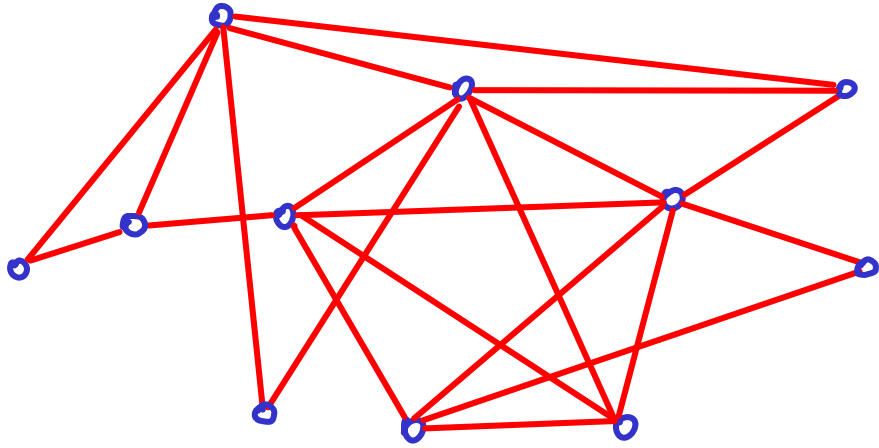
The induced subgraph obtained by removing all but  $S$  from  $V(G)$  is a complete graph ( $K_S$ )

# Independent Sets



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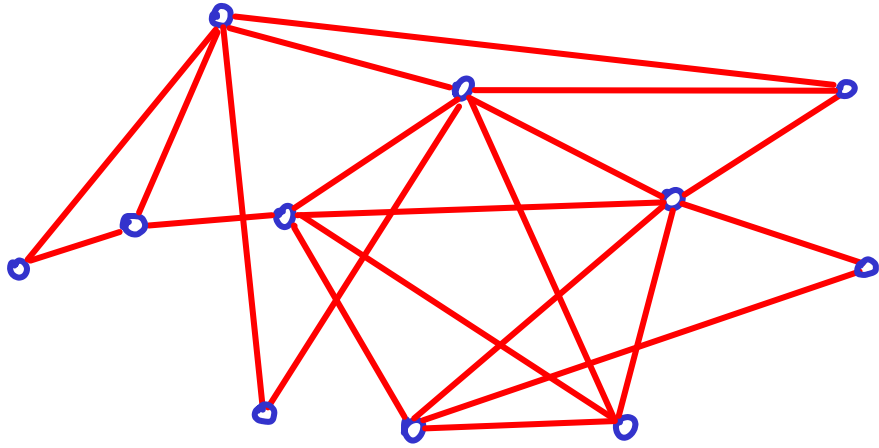
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# Independent Sets



Largest independent set?

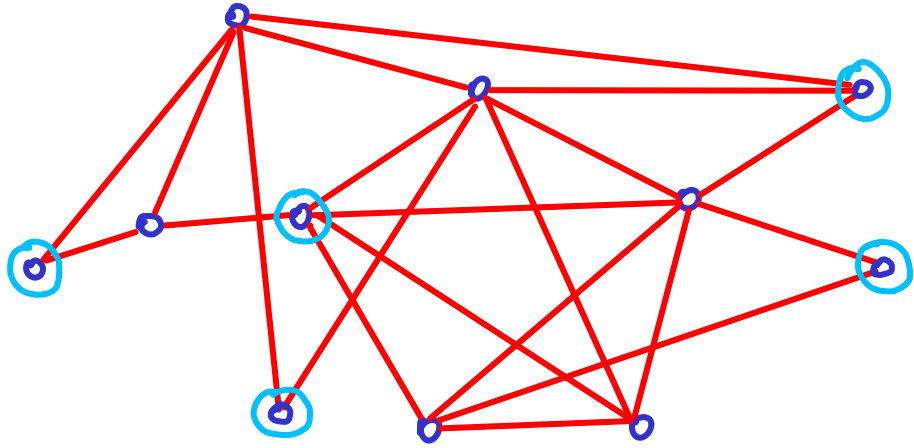
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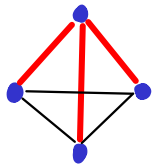




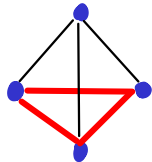
# Independent Sets



Largest independent set?



complements

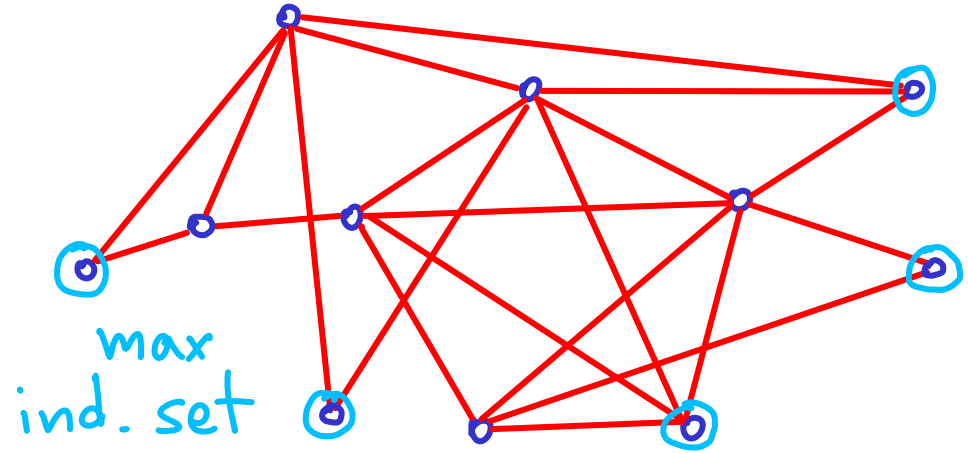
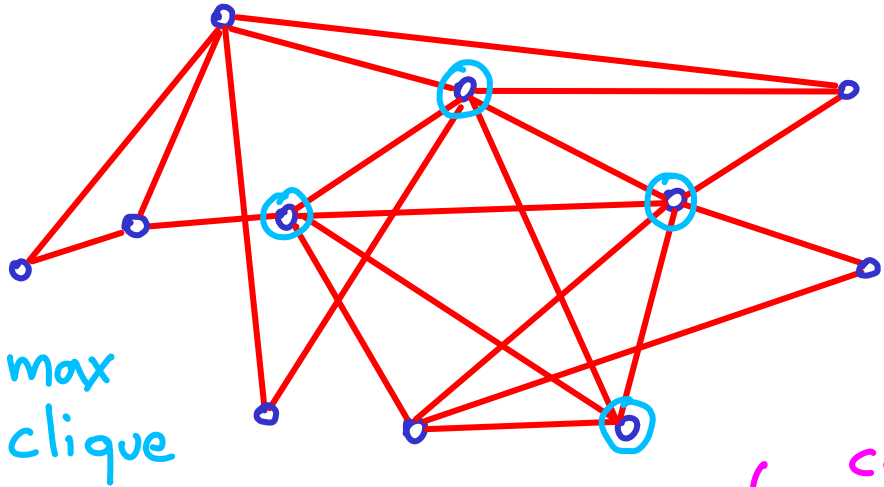


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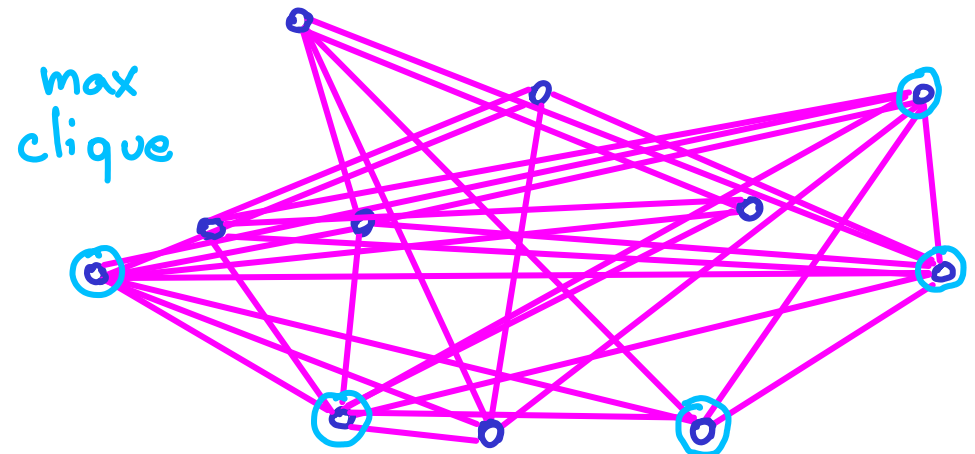
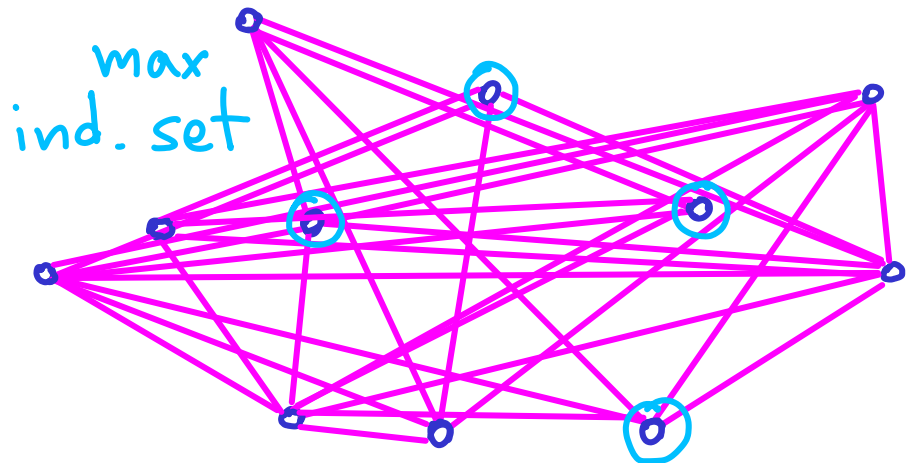
The induced subgraph obtained by removing all but  $S$  from  $V(G)$  is an edgeless graph.

↪ i.e. its complement is a complete graph

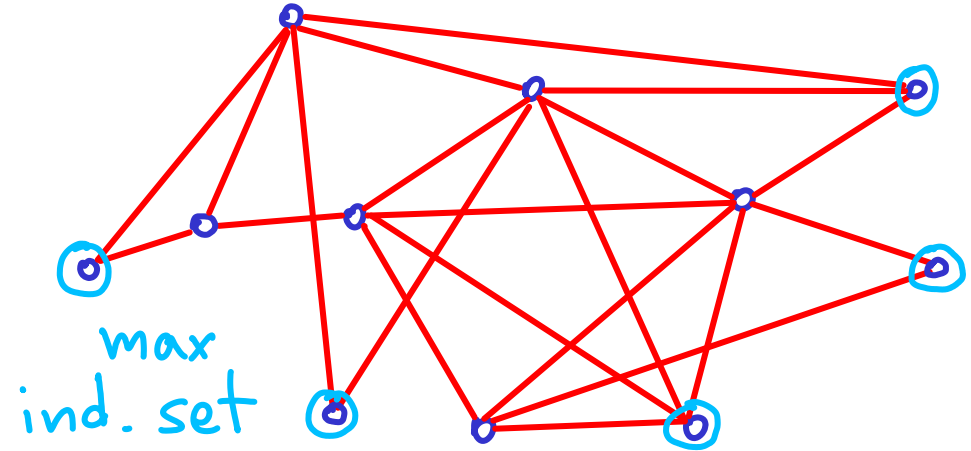
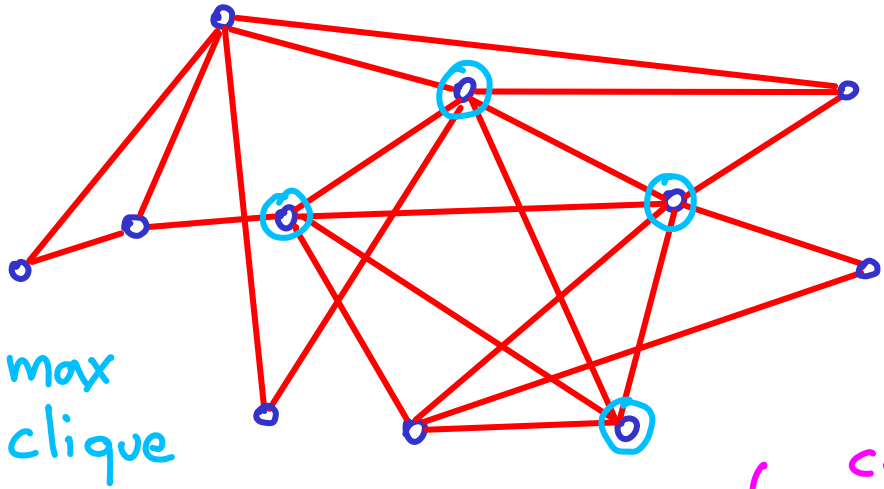
# cliques vs. independent sets



complement

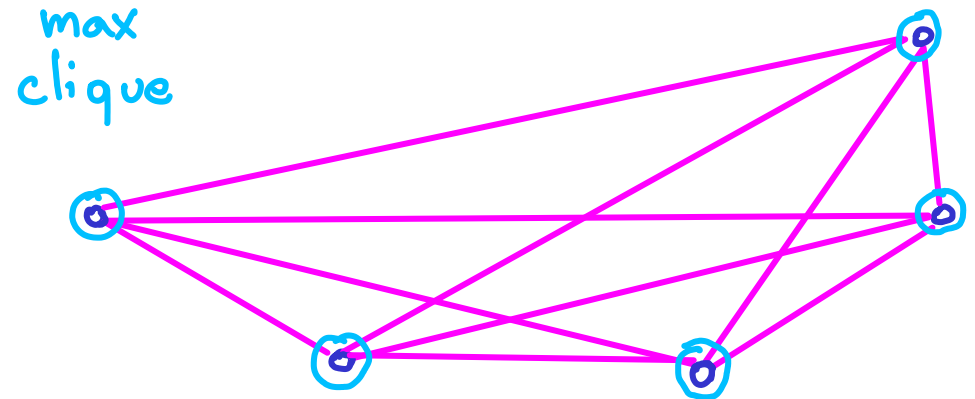


# cliques vs. independent sets



complement

A purple bracket under the word "complement" indicates that the highlighted nodes in the two graphs above are complementary sets.



Claim: Every graph with  $|V| \geq 6$  contains  
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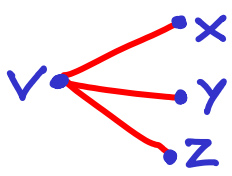
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If  $\underline{d(v) \geq 3}$  we have  $v$   } If  $\widehat{xy}$  or  $\widehat{xz}$  or  $\widehat{yz}$ : ?

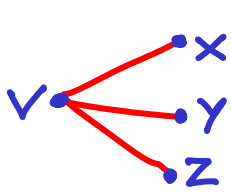


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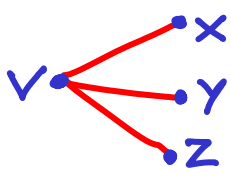
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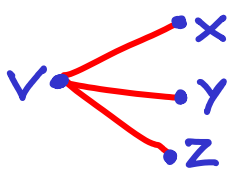
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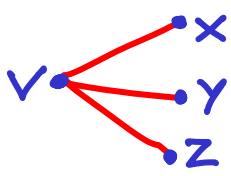
If  $d(v) \leq 2$  ... ?

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If  $\widehat{ab}, \widehat{bc}, \widehat{ac}$  are edges ... ?

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Otherwise one edge is missing (w.l.o.g.  $ab$ ) ...

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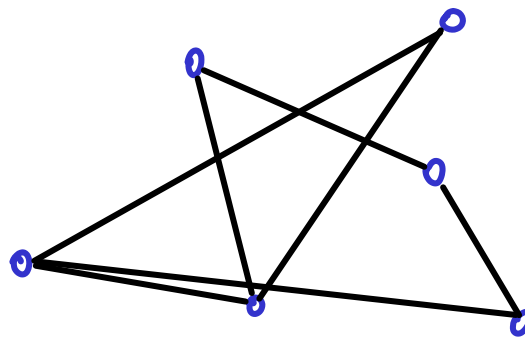
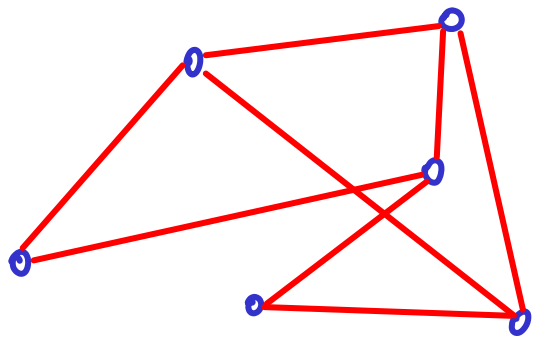
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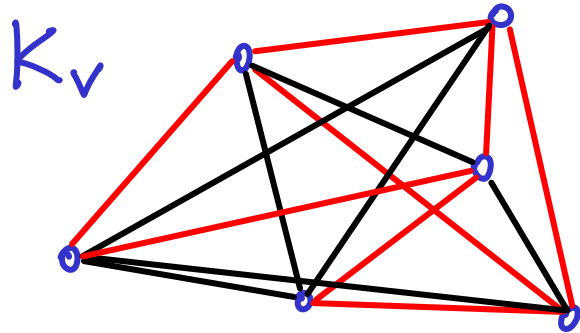
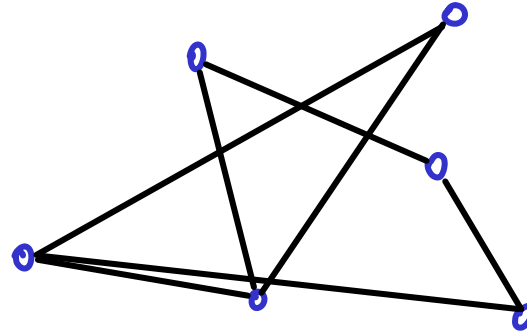
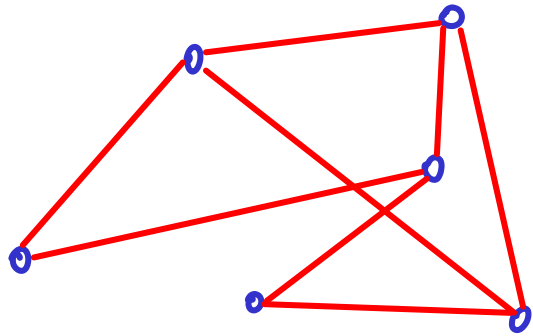
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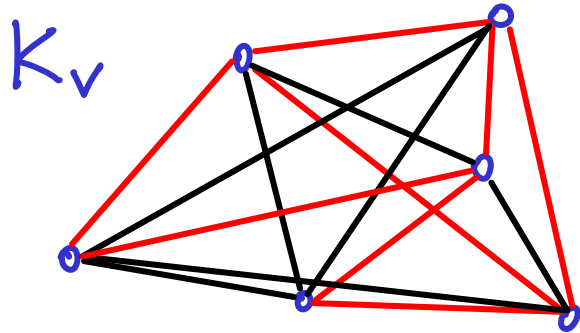
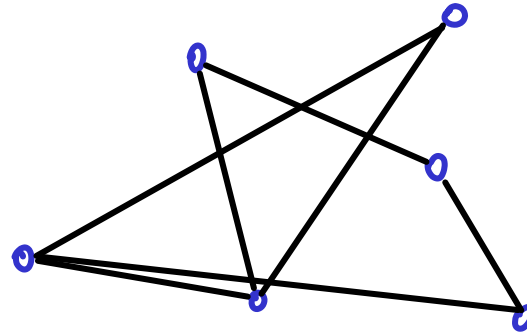
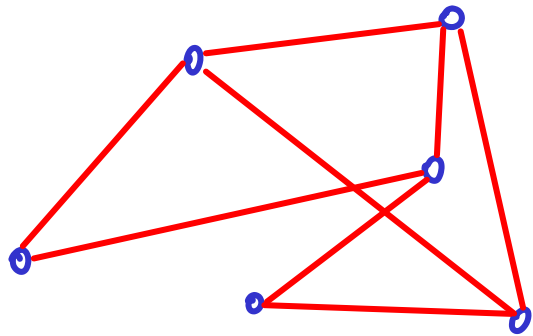
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Equivalent statement

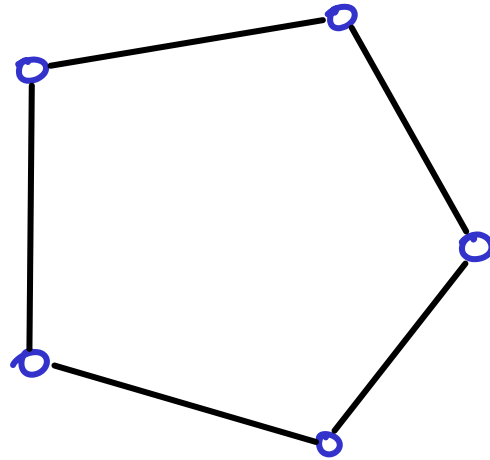
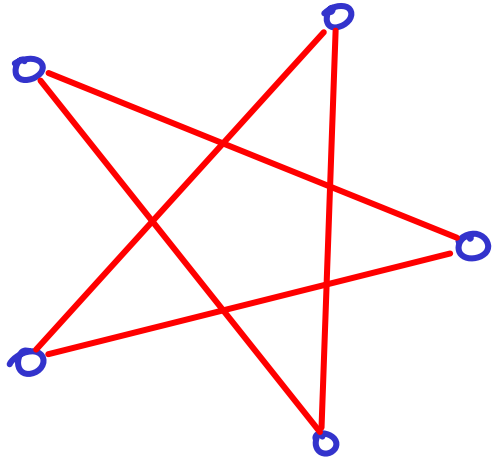
If we color each edge of  $K_v$  red or black, then we must get a red triangle or a black triangle

Recap: if you want a clique or an independent set of size 3  
then you'll be happy as long as  $|V| \geq \underline{6}$

$$|V| < 6 \quad ?$$

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$$R(4,4)$$



18

(one direction to be shown)



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18

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we don't know!  
[43...49]

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~ exponential

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even for small values  
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For more, see Ramsey's Thm.

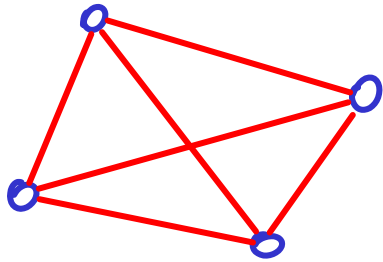
$R(x, y)$  : smallest number  $N$  such that any graph with  $\geq N$  vertices has a clique of size  $x$  or an independent set of size  $y$

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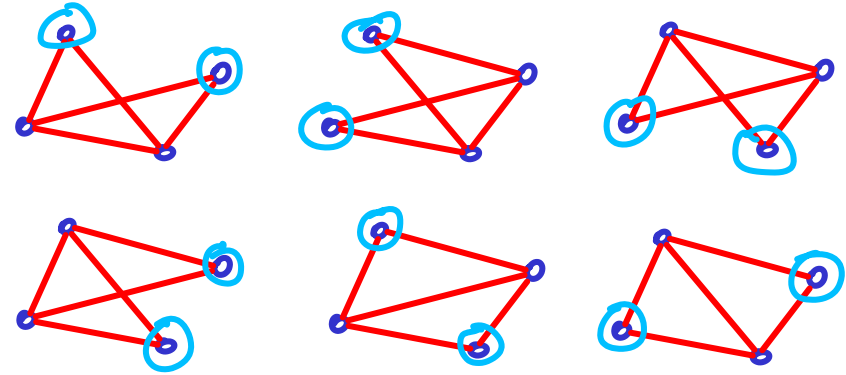
$$R(4, 2) = ?$$

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$$R(4,2) = 4$$



OR



$R(4,3)$

$R(4,3)$

Suppose  $|V| \geq 10$

Pick any vertex,  $A$ .  $\geq 9$  vertices remain.

$A$   
•

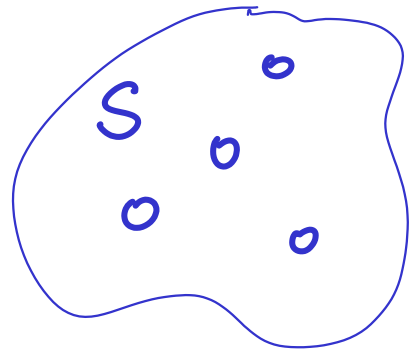


$R(4,3)$

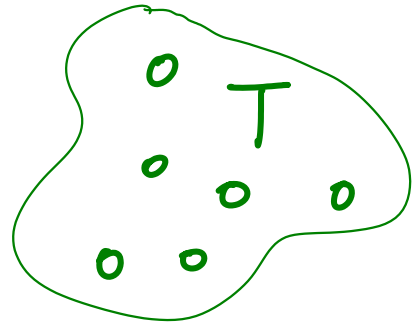
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Form 2 groups:  
 $S$  &  $T$



$A$

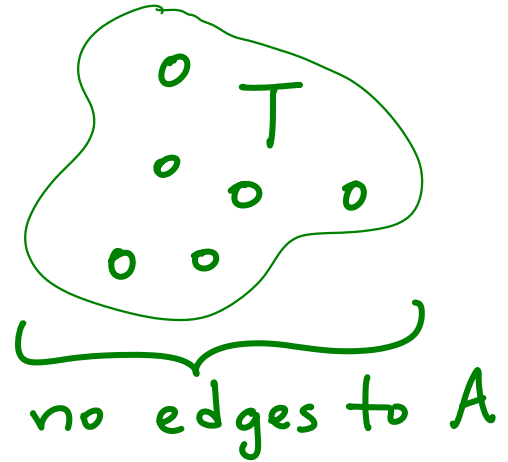
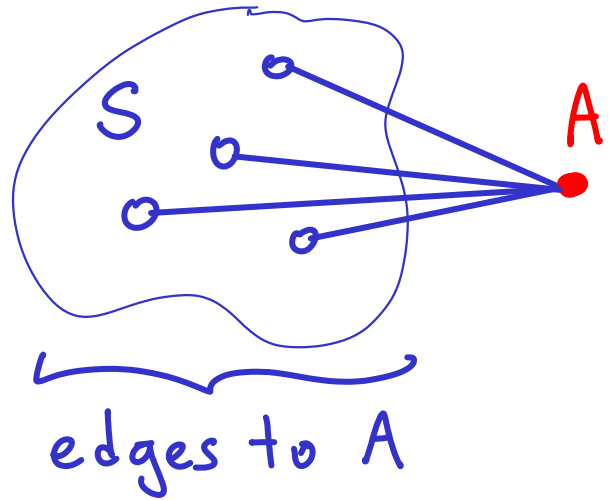


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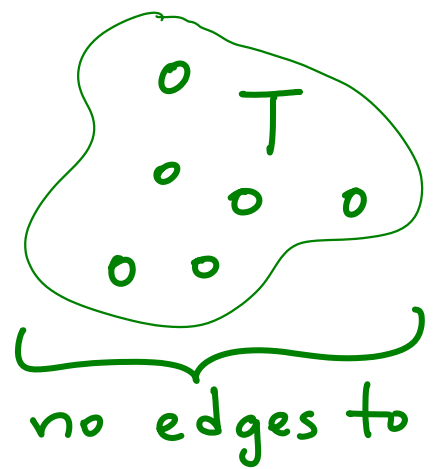
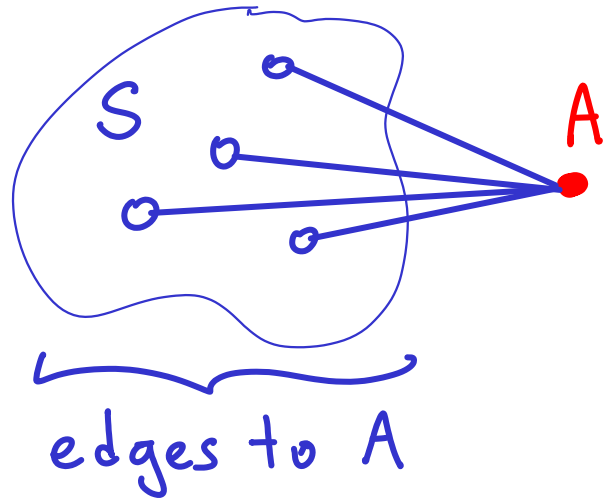


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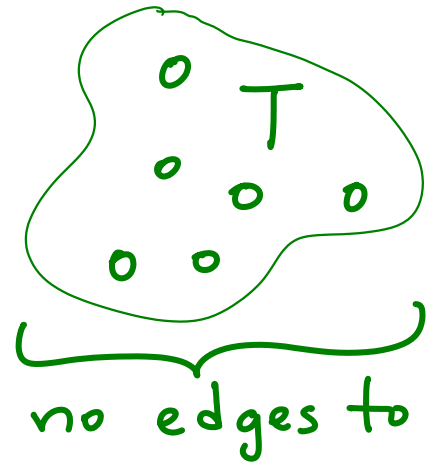
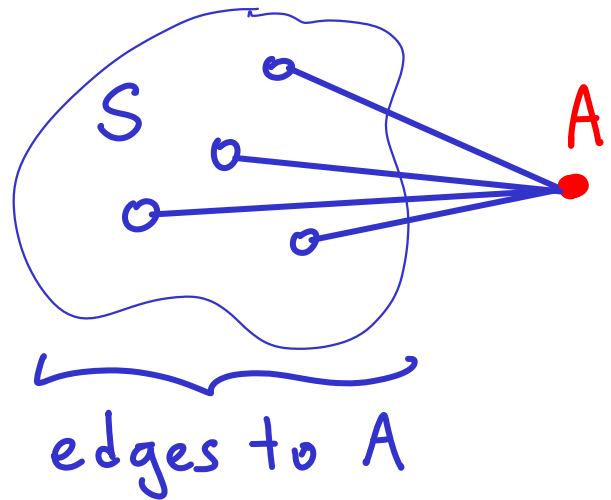
} either  
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 or  
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$R(4,3)$

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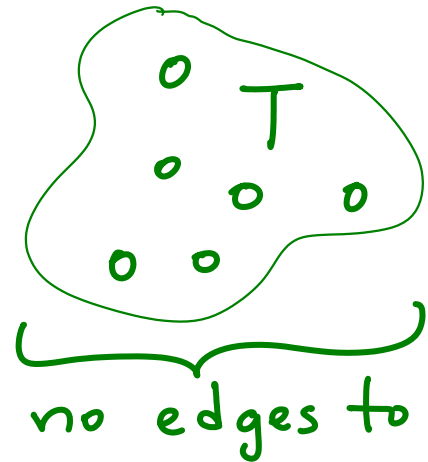
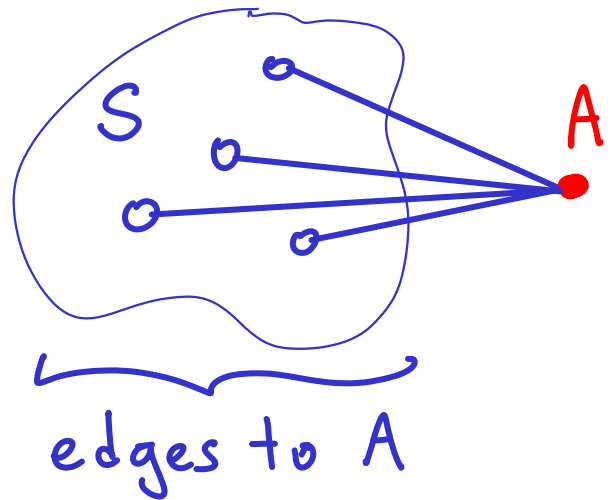
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$R(4,3)$

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 $S$  &  $T$



} either  
 $|S| \geq 6$   
or  
 $|T| \geq 4$

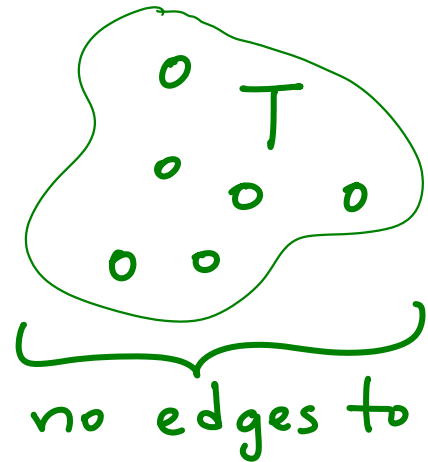
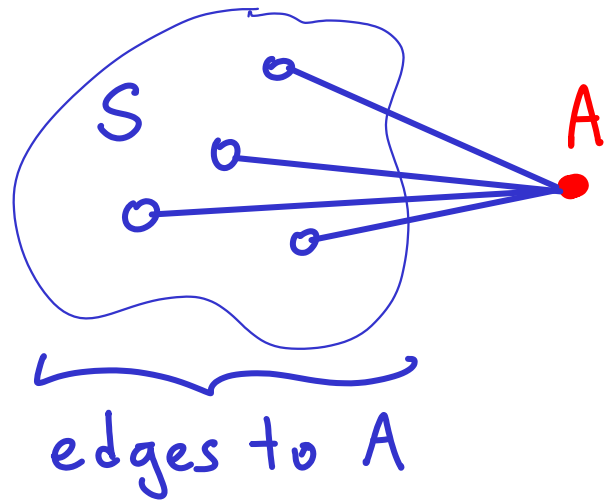
If  $|S| \geq 6$ , use  $R(3,3) = 6 \dots ?$

$R(4,3)$

Suppose  $|V| \geq 10$

Pick any vertex,  $A$ .  $\geq 9$  vertices remain.

Form 2 groups:  
 $S$  &  $T$



} either  
 $|S| \geq 6$   
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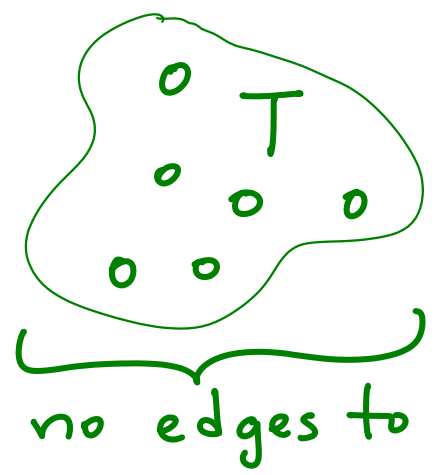
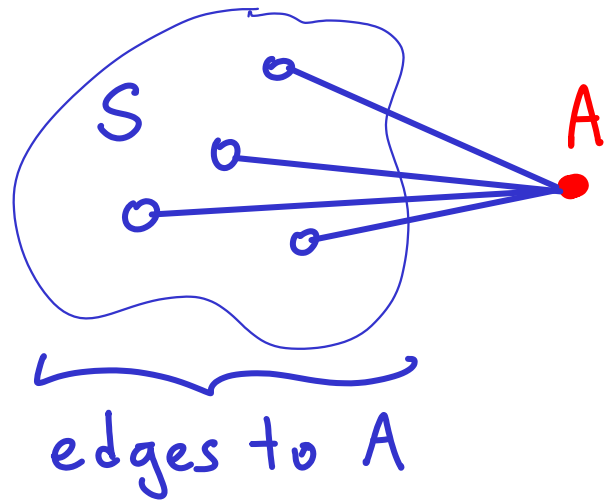
If  $|S| \geq 6$ , use  $R(3,3) = 6$  :  $S$  has 3 independent vertices  
OR  $S$  has a 3-clique  
...?

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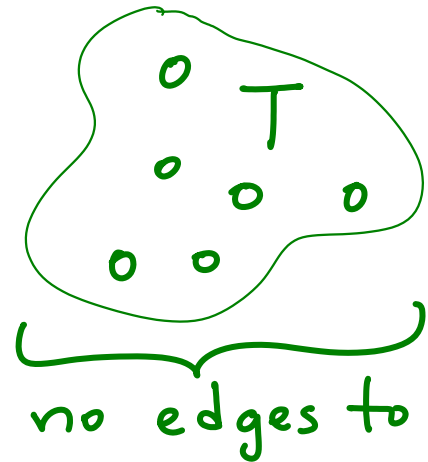
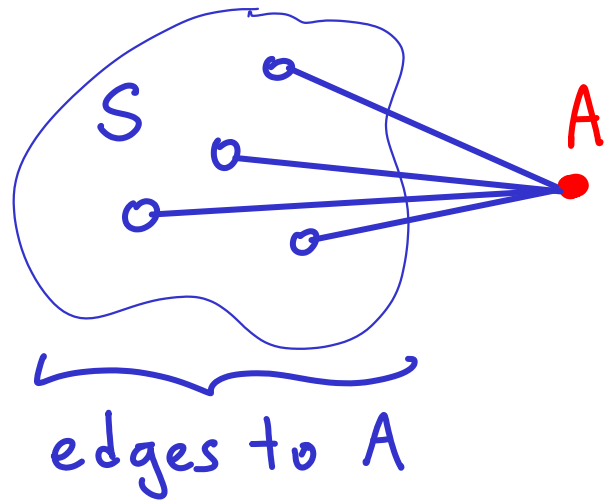
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Form 2 groups:  
 $S$  &  $T$



} either  
 $|S| \geq 6$   
or  
 $|T| \geq 4$

no edges to A

If  $|S| \geq 6$ , use  $R(3,3) = 6$  :  $S$  has 3 independent vertices (done),  
OR  $S$  has a 3-clique, so with  $A$  we get a 4-clique.

If  $|T| \geq 4$  ?

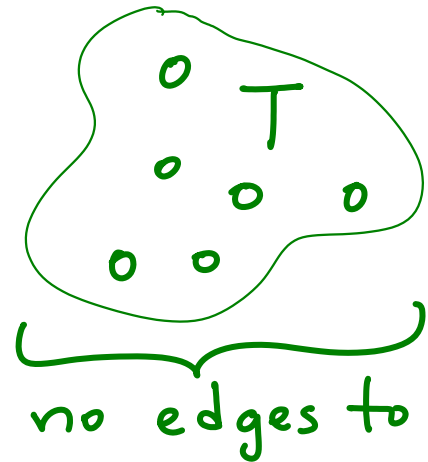
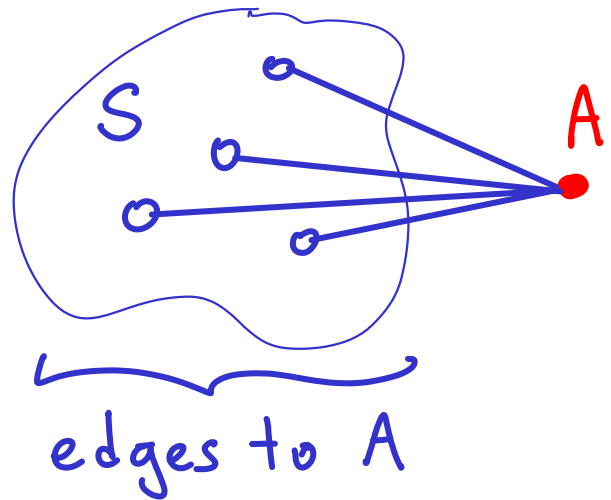


$R(4,3)$

Suppose  $|V| \geq 10$

Pick any vertex,  $A$ .  $\geq 9$  vertices remain.

Form 2 groups:  
 $S$  &  $T$



} either  
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If  $|S| \geq 6$ , use  $R(3,3) = 6$  :  $S$  has 3 independent vertices (done),  
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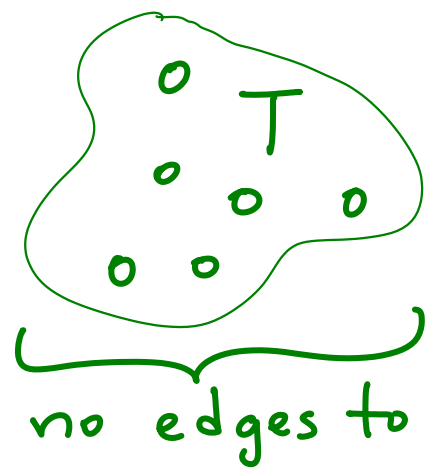
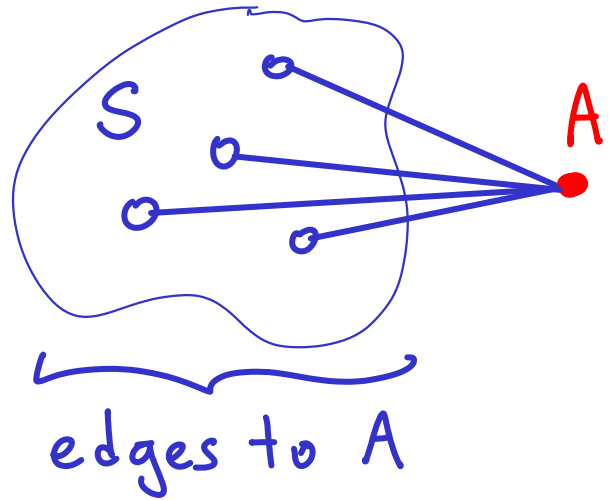
If  $|T| \geq 4$ , if  $T$  is a clique : done.

$R(4,3)$

Suppose  $|V| \geq 10$

Pick any vertex,  $A$ .  $\geq 9$  vertices remain.

Form 2 groups:  
 $S$  &  $T$



} either  
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If  $|S| \geq 6$ , use  $R(3,3) = 6$  :  $S$  has 3 independent vertices (done),  
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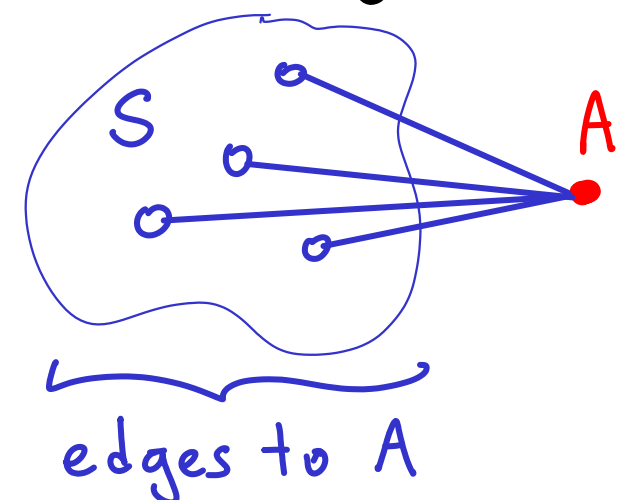
If  $|T| \geq 4$ , if  $T$  is a clique : done.  
Otherwise  $\exists a, b$  in  $T$  w/ no edge. Combine w/  $A$ .

$R(4,3) \leq 10$

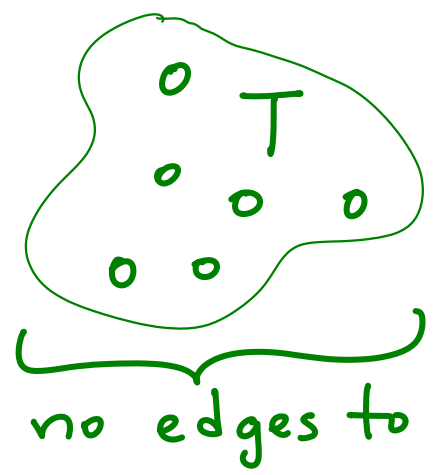
Suppose  $|V| \geq 10$

Pick any vertex,  $A$ .  $\geq 9$  vertices remain.

Form 2 groups:  
 $S$  &  $T$



edges to  $A$



no edges to  $A$

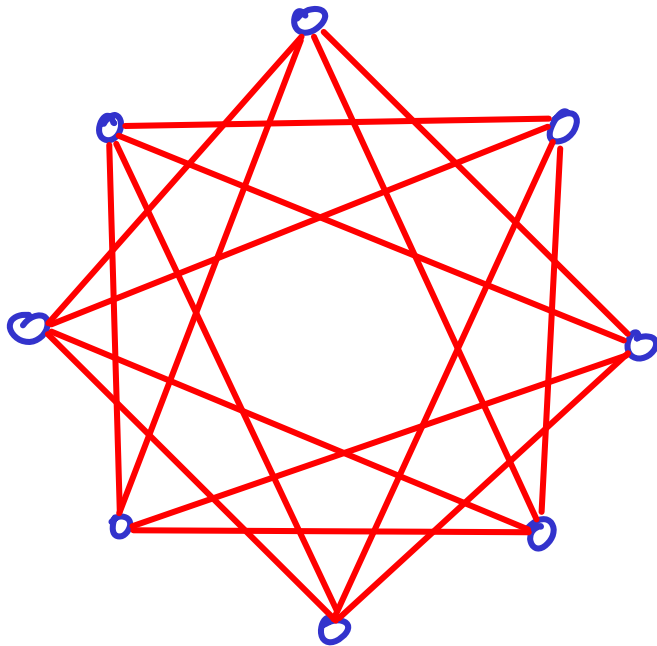
} either  
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If  $|S| \geq 6$ , use  $R(3,3) = 6$  :  $S$  has 3 independent vertices (done),  
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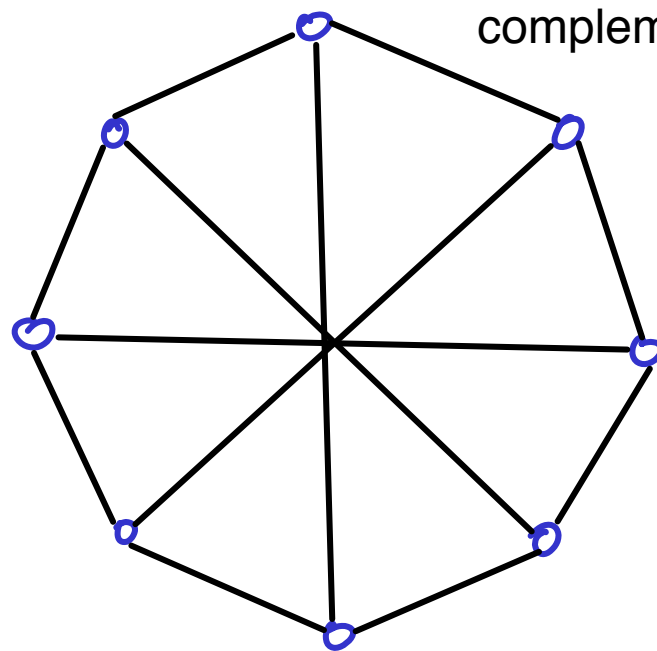
If  $|T| \geq 4$ , if  $T$  is a clique : done.  
Otherwise  $\exists a, b$  in  $T$  w/ no edge. Combine w/  $A$ .  $\square$

$$R(4,3) \leq 10$$

$$\underline{R(4,3) > 8}$$



no 4-clique



complement(G)

no 3-clique = no 3-independent in G

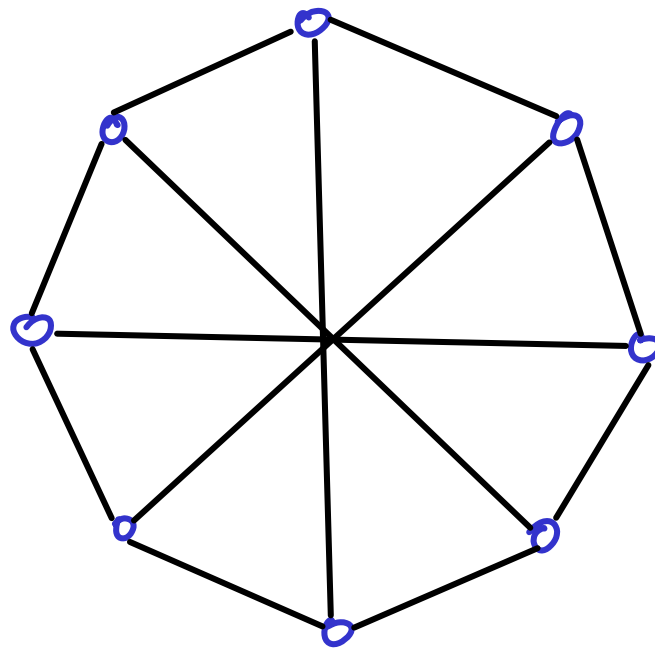
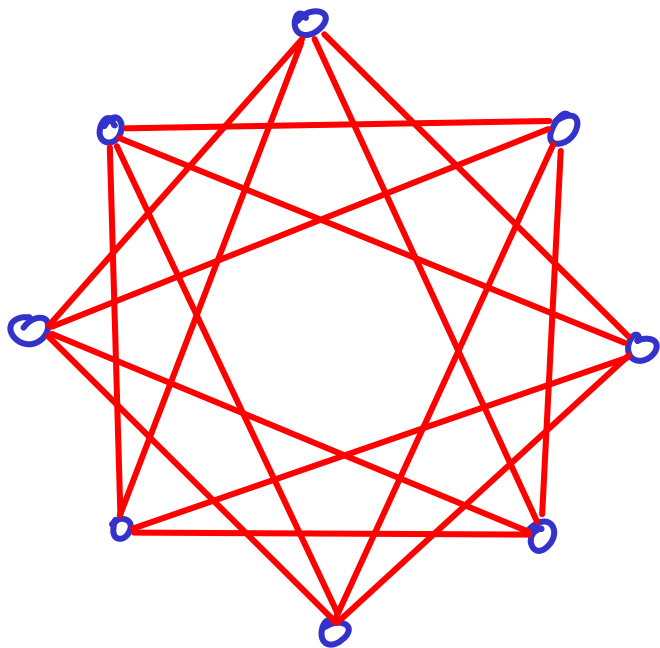
$$R(4,3) \leq 10$$

$$\underline{R(4,3) > 8}$$

... turns out  $R(4,3) = 9$

↳ not terribly hard

↳ notice  $R(x,y) = R(y,x)$



$R(4,4)$

$R(4,4)$

Suppose  $|V| \geq 18$

Pick any vertex,  $A$ .  $\geq 17$  vertices remain.

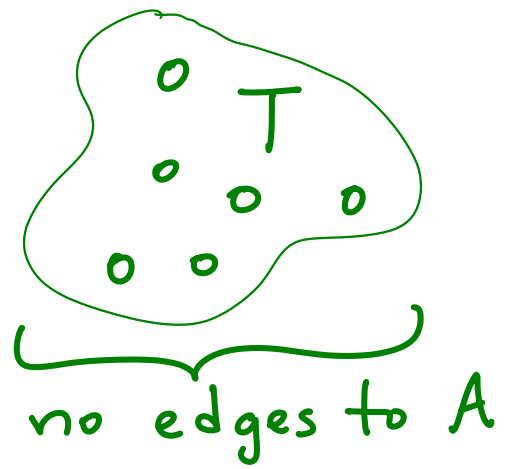
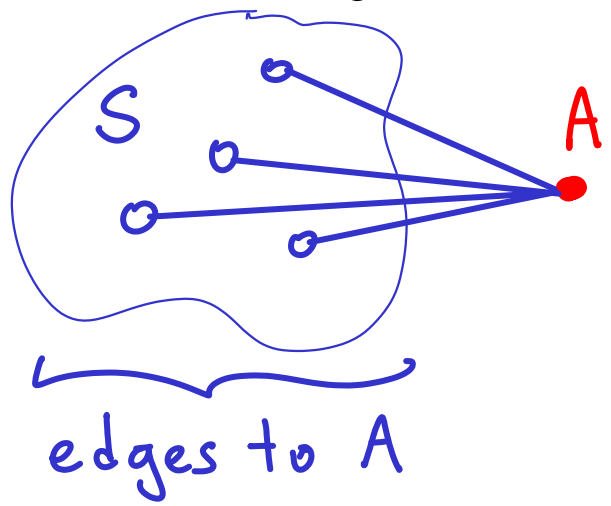
$A$   
•

$R(4,4)$

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Form 2 groups:  
 $S$  &  $T$



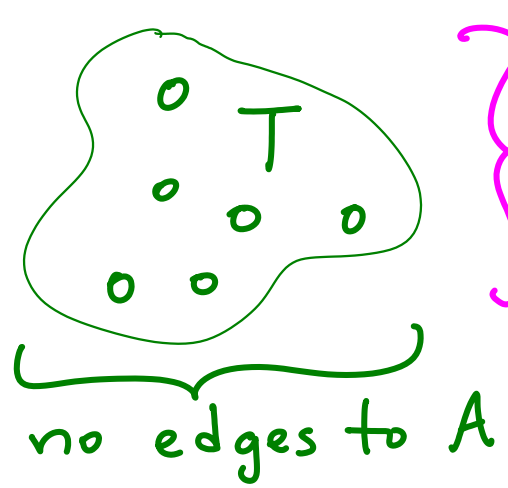
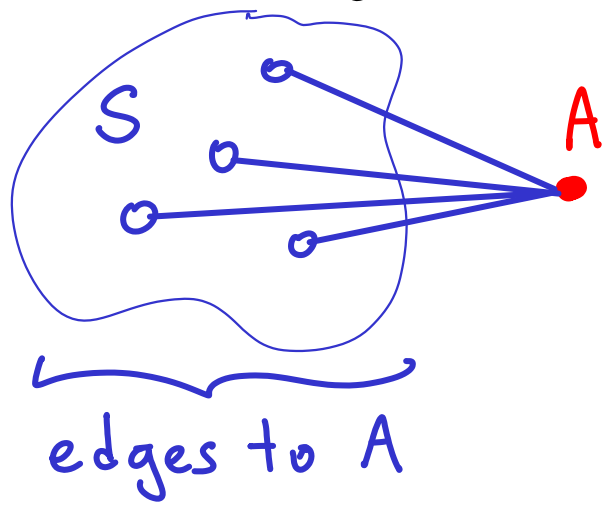


$R(4,4)$

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} either  
 $|S| \geq 9$   
or  
 $|T| \geq 9$

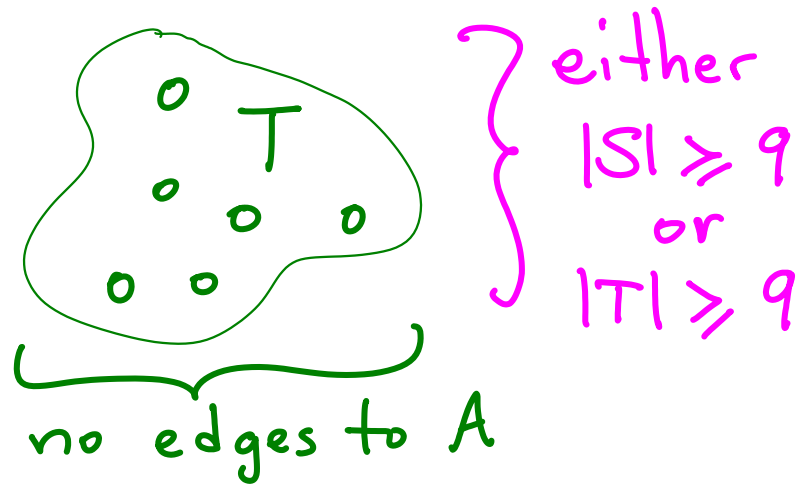
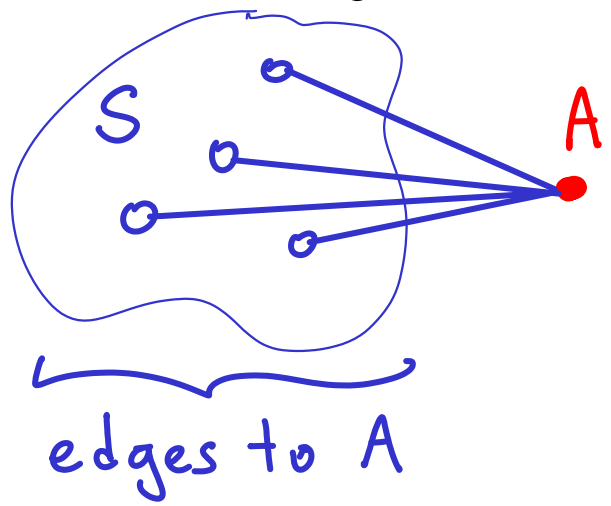
if  $|S| \geq 9$  ?

$R(4,4)$

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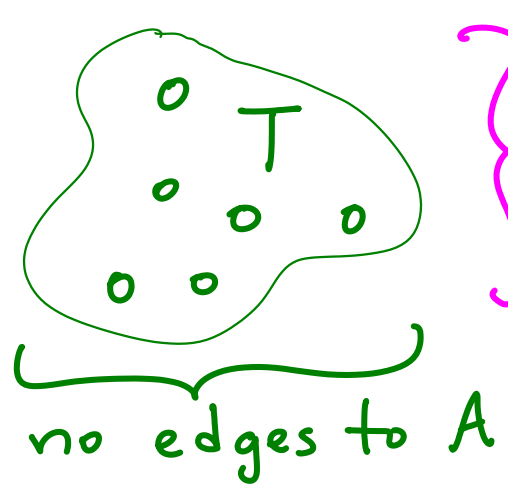
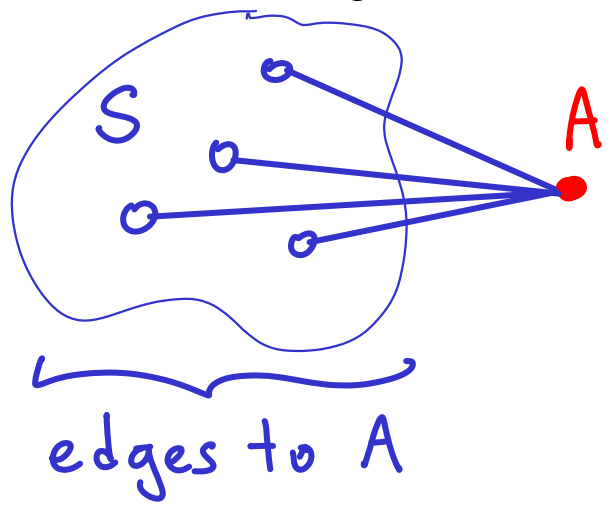
If  $|S| \geq 9$ , use  $R(3,4) = 9 \dots ?$

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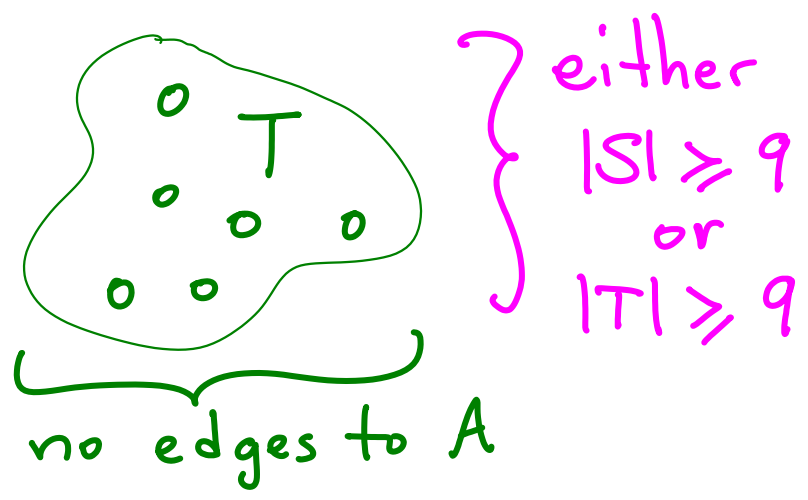
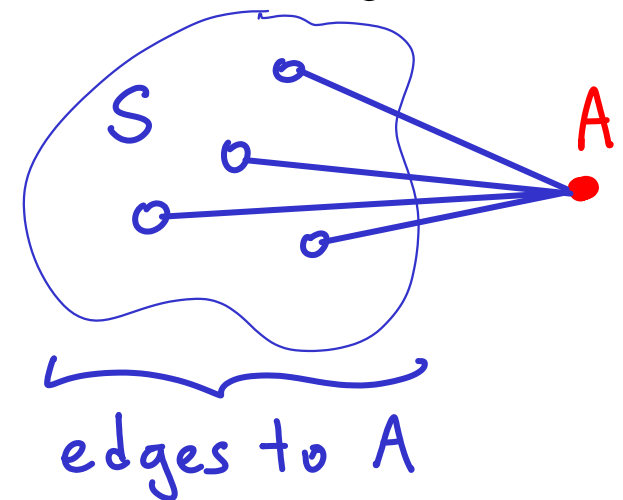
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OR  $S$  has a 3-clique ... ?

$R(4,4)$

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Form 2 groups:  
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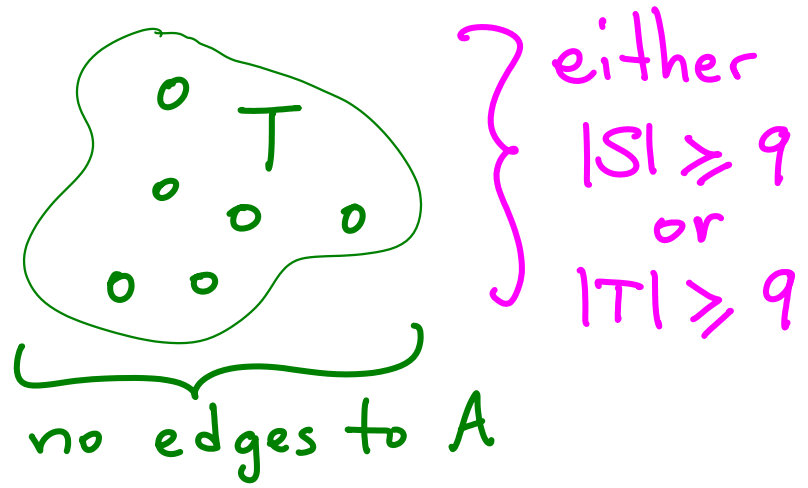
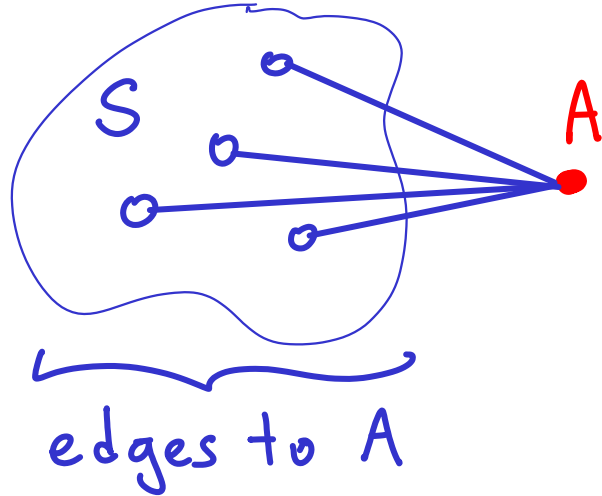
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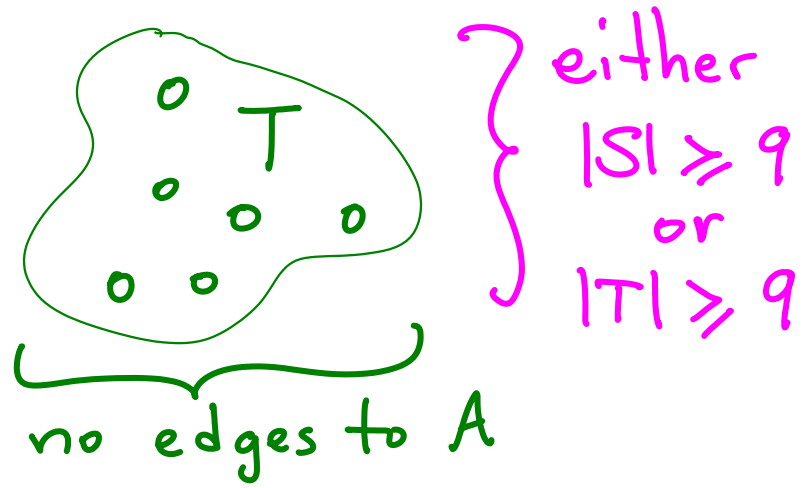
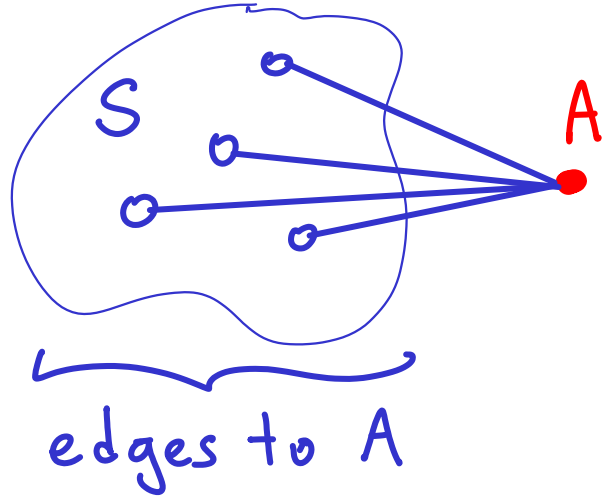
If  $|T| \geq 9$  ?

$R(4,4)$

Suppose  $|V| \geq 18$

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Form 2 groups:  
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If  $|S| \geq 9$ , use  $R(3,4) = 9$  :  $S$  has 4 independent vertices (done)  
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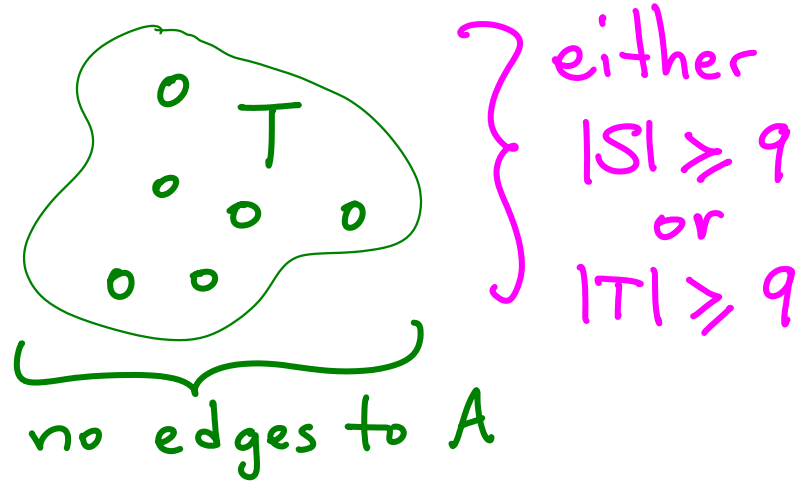
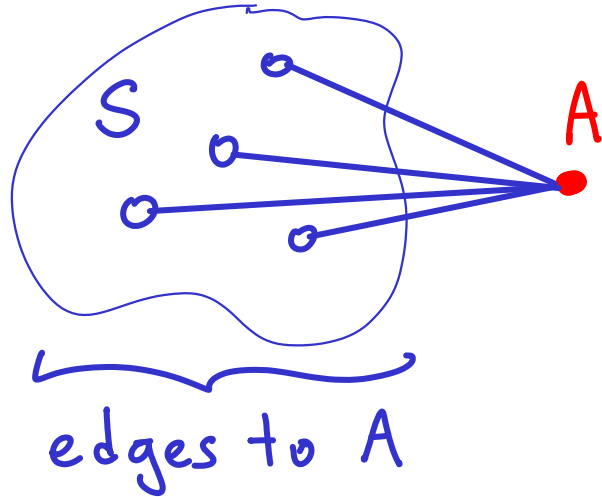
If  $|T| \geq 9$ , use  $R(4,3) = 9$  ... ?

$R(4,4)$

Suppose  $|V| \geq 18$

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If  $|S| \geq 9$ , use  $R(3,4) = 9$  :  $S$  has 4 independent vertices (done)  
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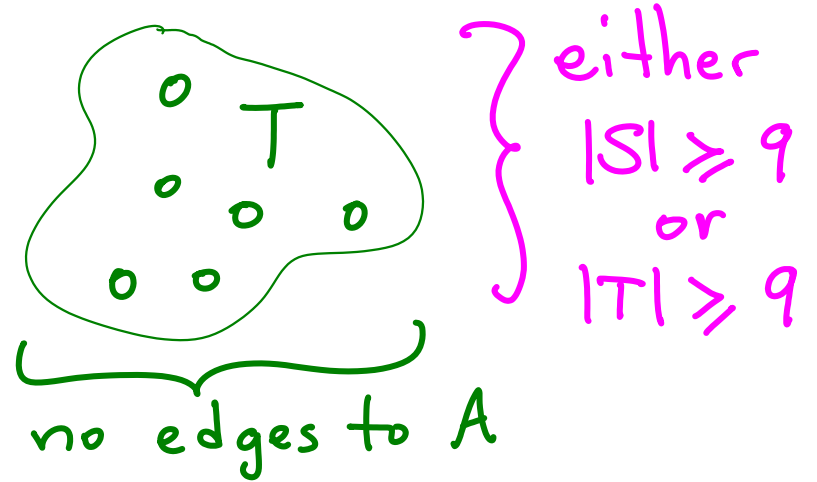
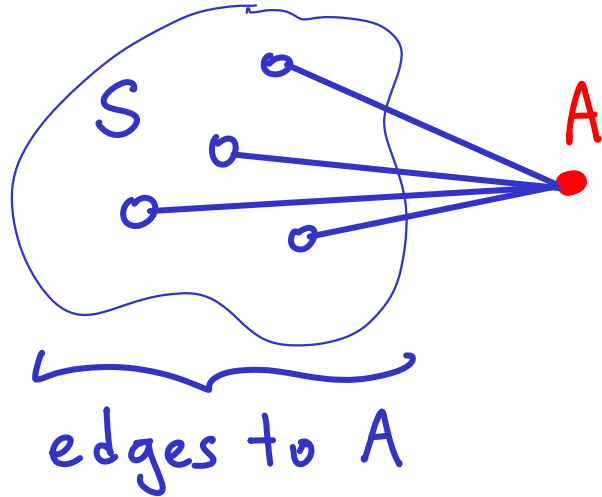
If  $|T| \geq 9$ , use  $R(4,3) = 9$  :  $T$  has a 4-clique, OR  
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$R(4,4)$

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If  $|S| \geq 9$ , use  $R(3,4) = 9$  :  $S$  has 4 independent vertices (done) OR  $S$  has a 3-clique, so with  $A$  we get a 4-clique.

If  $|T| \geq 9$ , use  $R(4,3) = 9$  :  $T$  has a 4-clique, (done) OR  $T$  has 3 independent vertices, so with  $A$  we have 4.

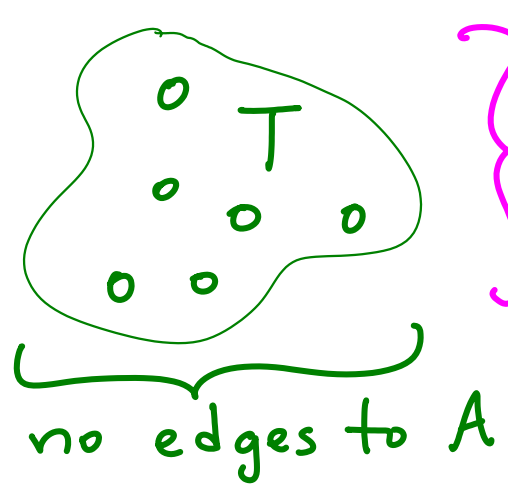
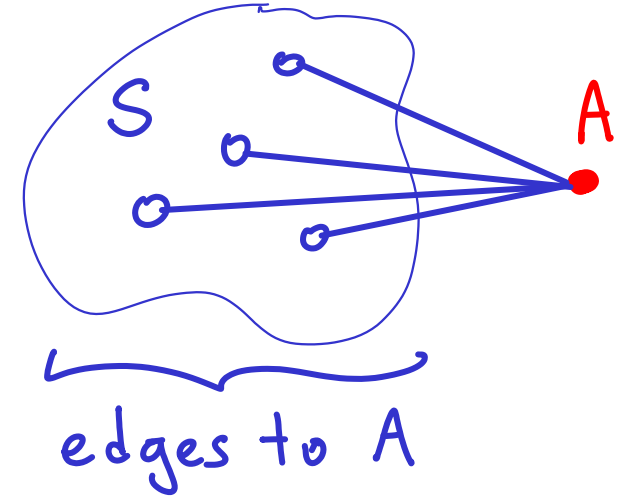


$R(4,4) \leq 18$

Suppose  $|V| \geq 18$

Pick any vertex,  $A$ .  $\geq 17$  vertices remain.

Form 2 groups:  
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} either  
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Notes:

(1) if we only knew that  $R(4,3) \leq 10$  (instead of  $= 9$ )

then ... ?

Notes:

(i) if we only knew that  $R(4,3) \leq 10$  (instead of  $= 9$ )  
we could have used  $|V| \geq 20$  for  $R(4,4)$

As you bound smaller  $R(\ )$  values, you can get (loose) bounds for larger ones

# Notes:

(1) if we only knew that  $R(4,3) \leq 10$  (instead of  $=9$ )  
we could have used  $|V| \geq 20$  for  $R(4,4)$

(2) there is a graph w/ 17 vertices with no  
clique or independent set of size 4

$$\hookrightarrow R(4,4) = 18$$