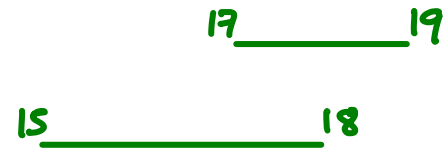
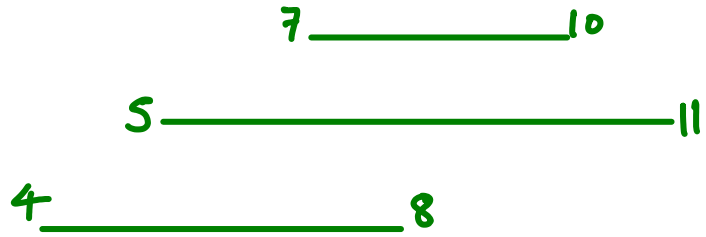


# INTERVAL TREES

set  $S$   
of  
intervals



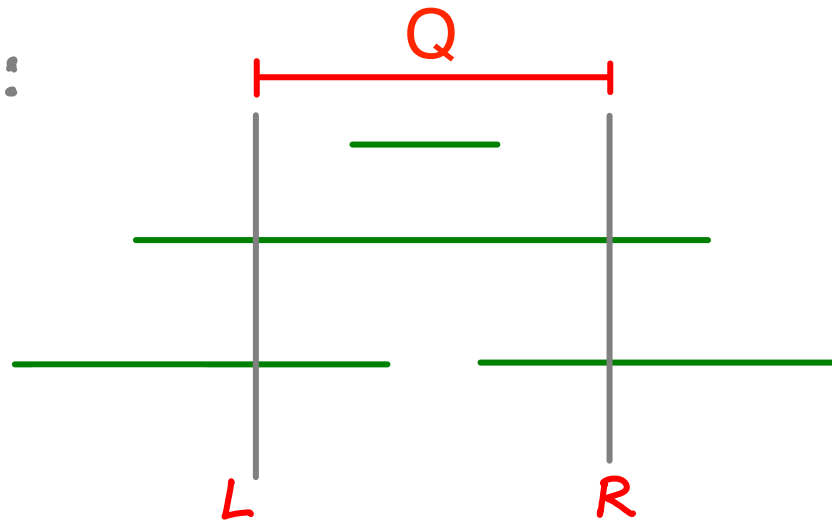
# INTERVAL TREES

set  $S$   
of  
intervals



Query : given an interval  $Q$ ,  
return any interval in the set  $S$  that partially overlaps  $Q$   
(if one exists)

types of overlap:



---

For  $Q$  to not overlap an input interval,  $Q$  must be:

entirely to the right

OR

entirely to the left

$high < L$

$R < low$

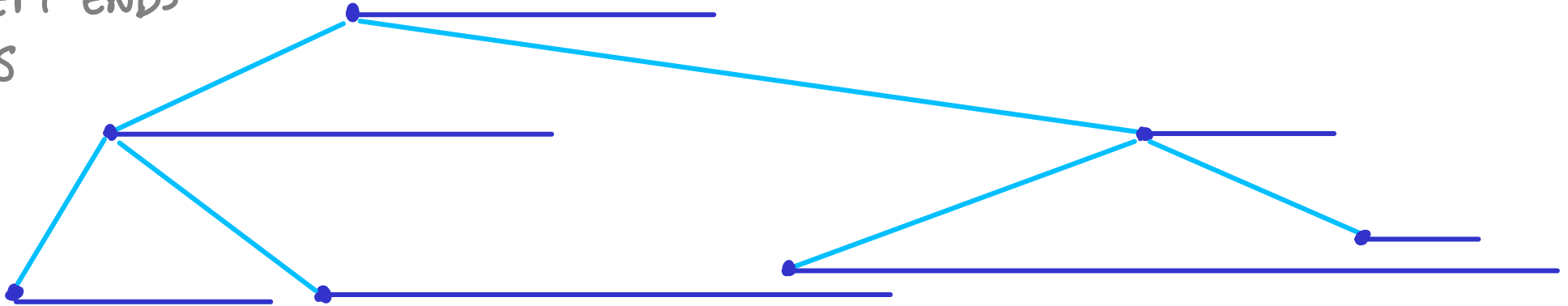


$O(1)$  time

1D:



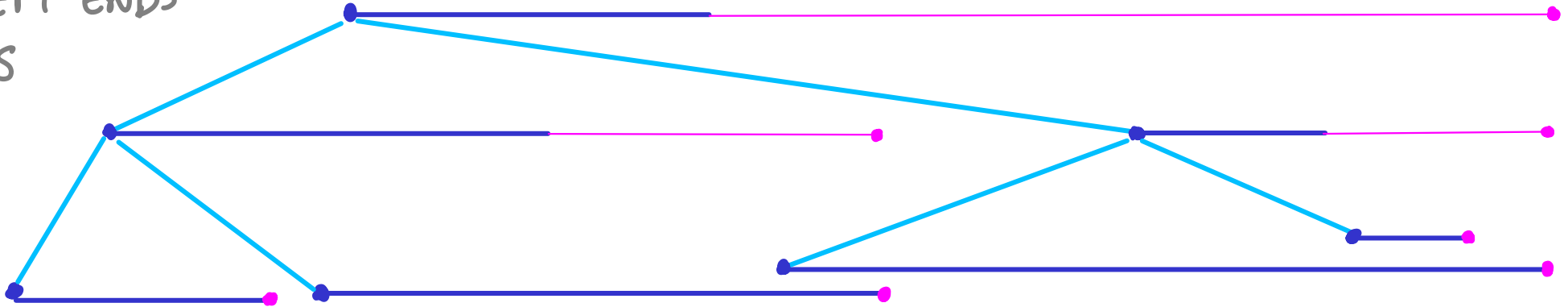
BST w/ LEFT ENDS  
as KEYS



1D:



BST w/ LEFT ENDS  
as KEYS



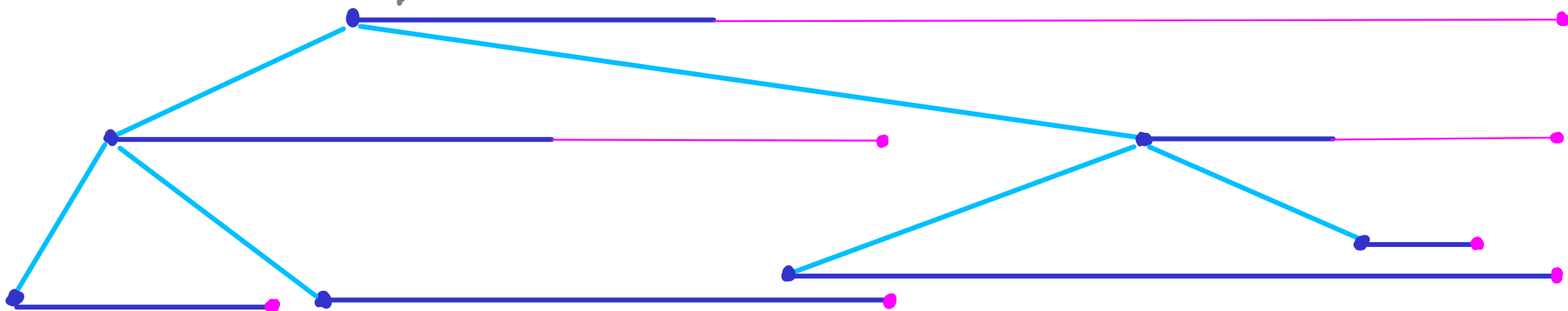
1D:



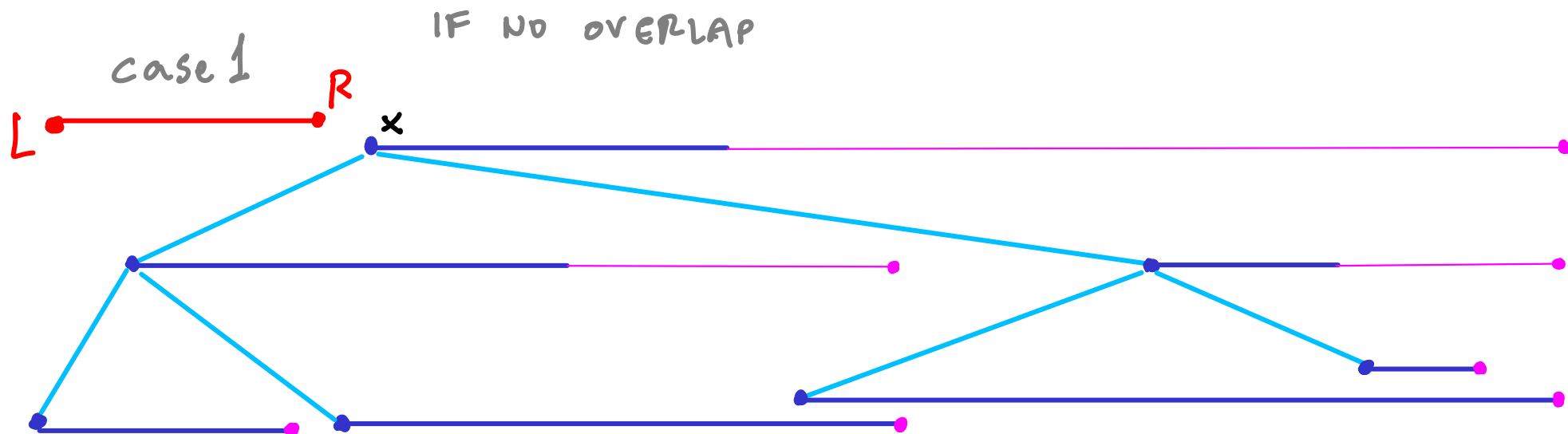
compare w/  
root first



If overlap, DONE: report root

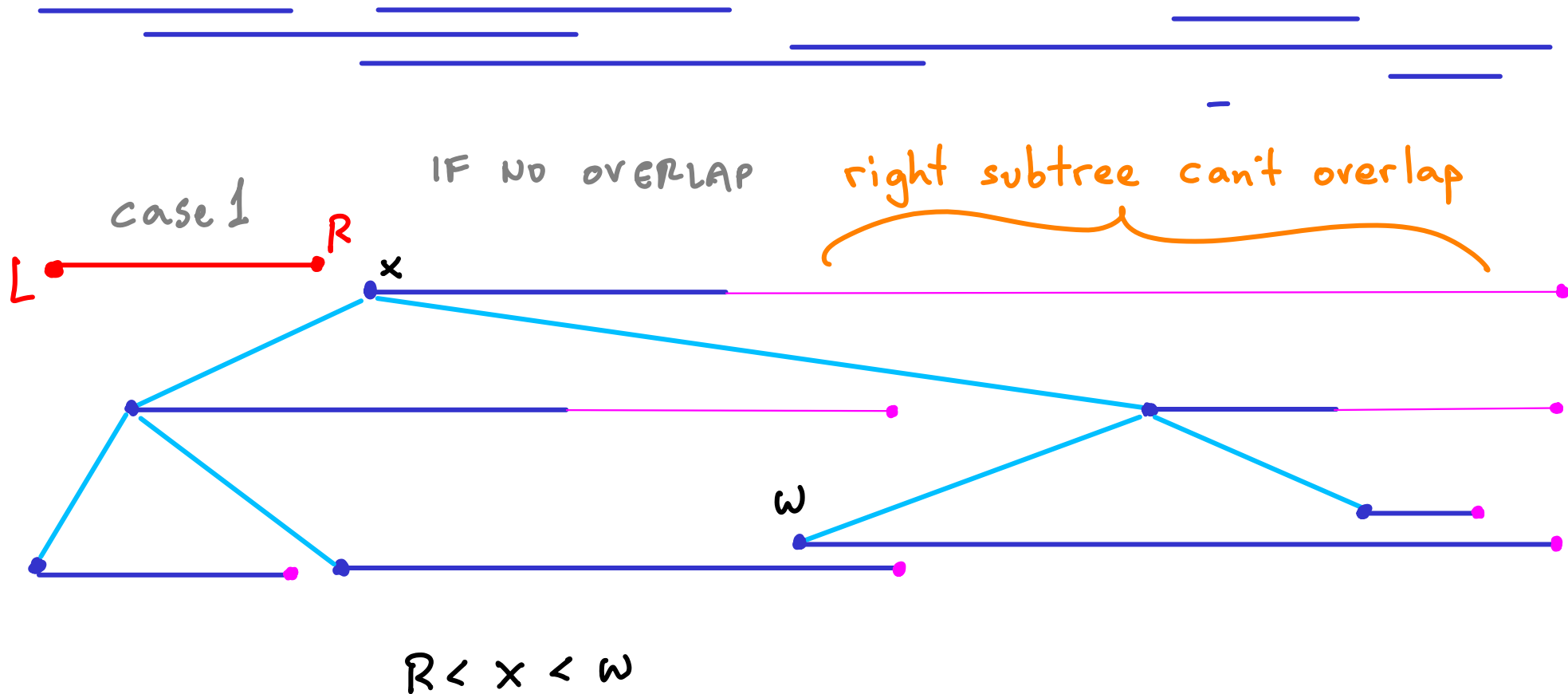


1D:



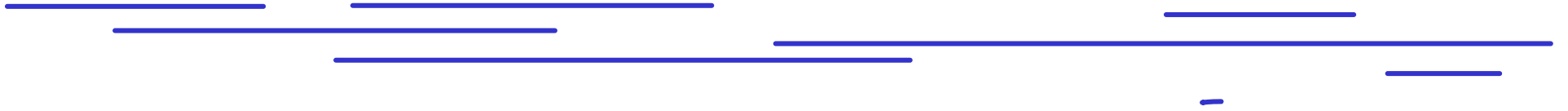
$$R < x$$

1D:





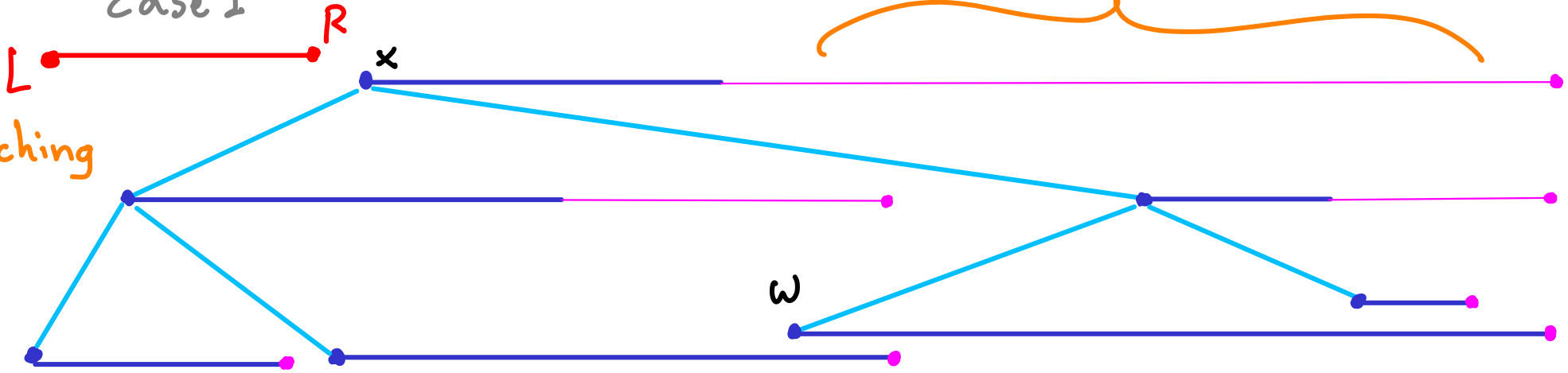
1D:



case 1

IF NO OVERLAP

right subtree can't overlap



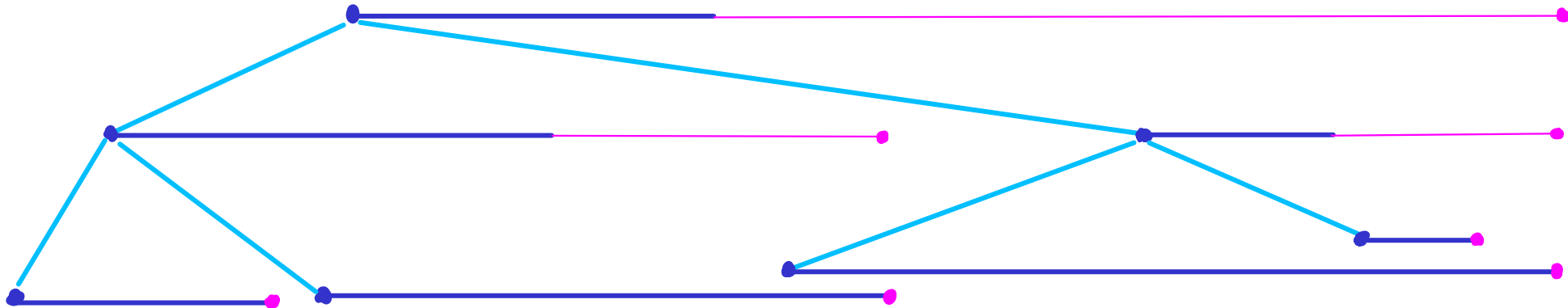
keep searching  
LEFT

$$R < x < w$$

1D:



IF NO OVERLAP



1D:

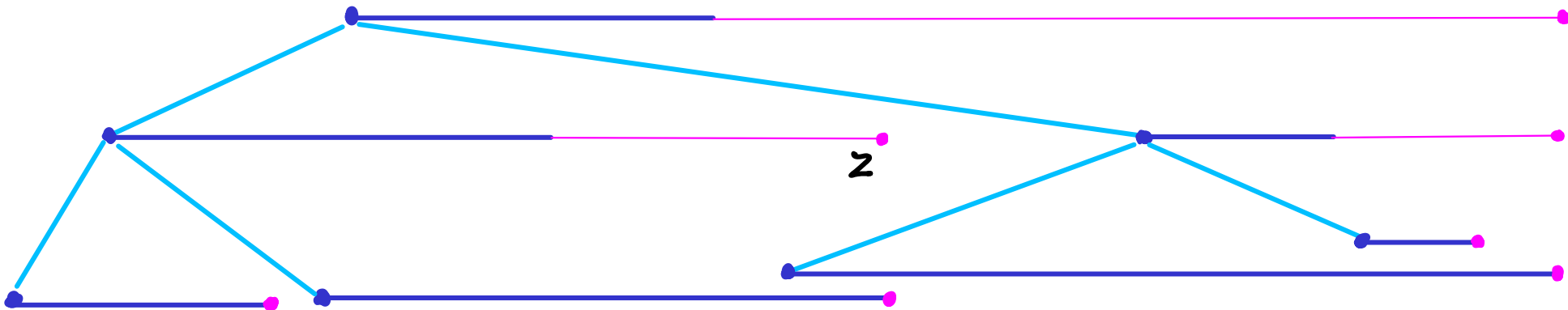


IF NO OVERLAP



IF  $z \geq L$

?



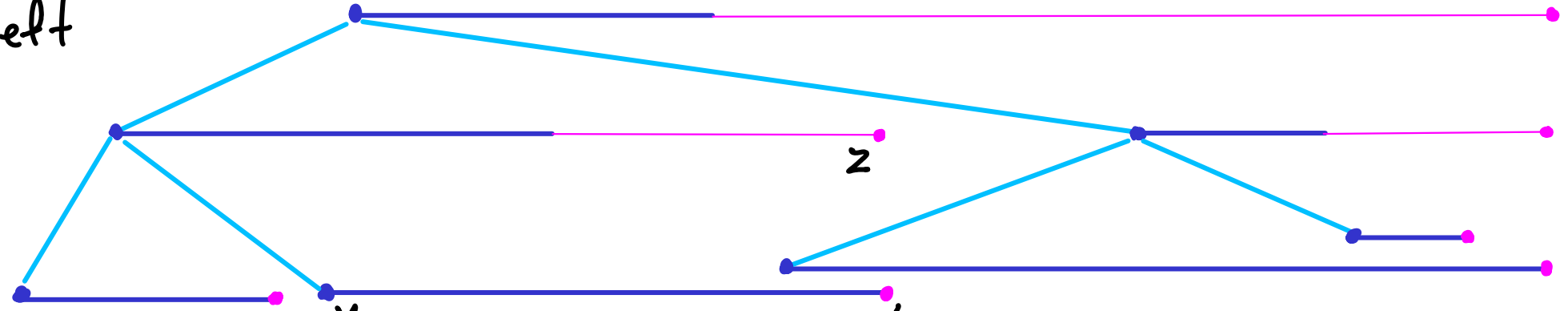
1D:



IF NO OVERLAP



IF  $z \geq L$   
 search left



$\exists y, z'$   
 s.t.  
 $y < L < z'$

} guaranteed overlap

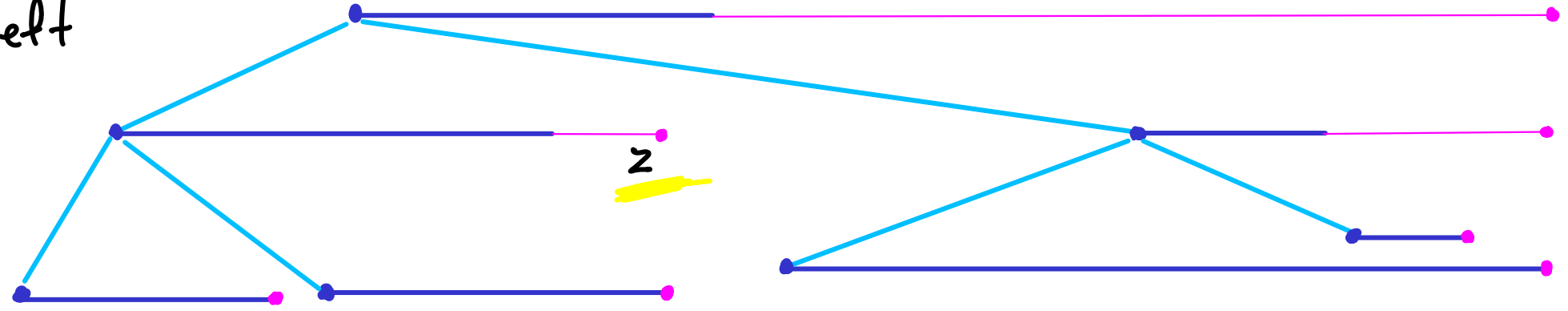
1D:



IF NO OVERLAP



IF  $z \geq L$   
search left



$\exists y, z'$   
 s.t.  $y < L < z'$

} guaranteed overlap

else ( $z < L$ )

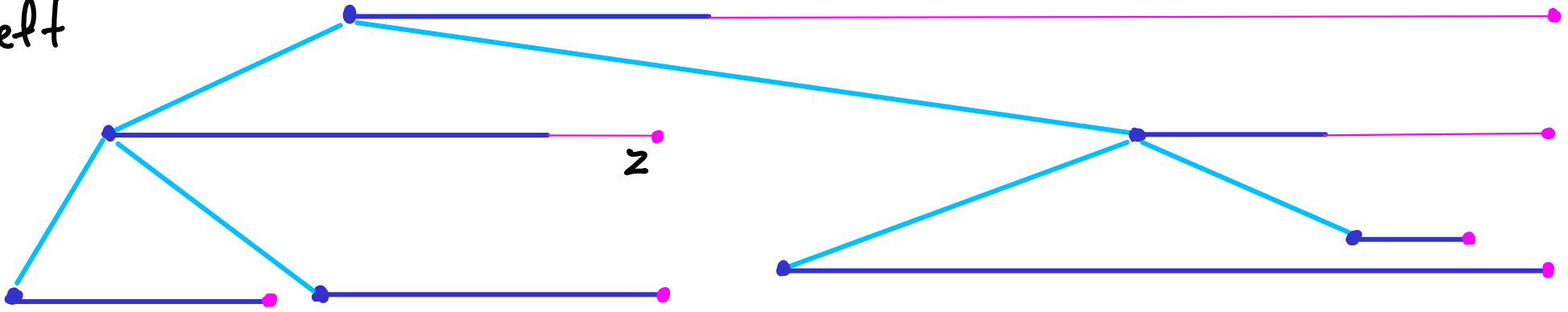
1D:



IF NO OVERLAP



IF  $z \geq L$   
 search left



$\exists \overline{yz'}$   
 s.t.  $y < L < z'$

} guaranteed overlap

else ( $z < L$ )

- NO overlap to left
- search right

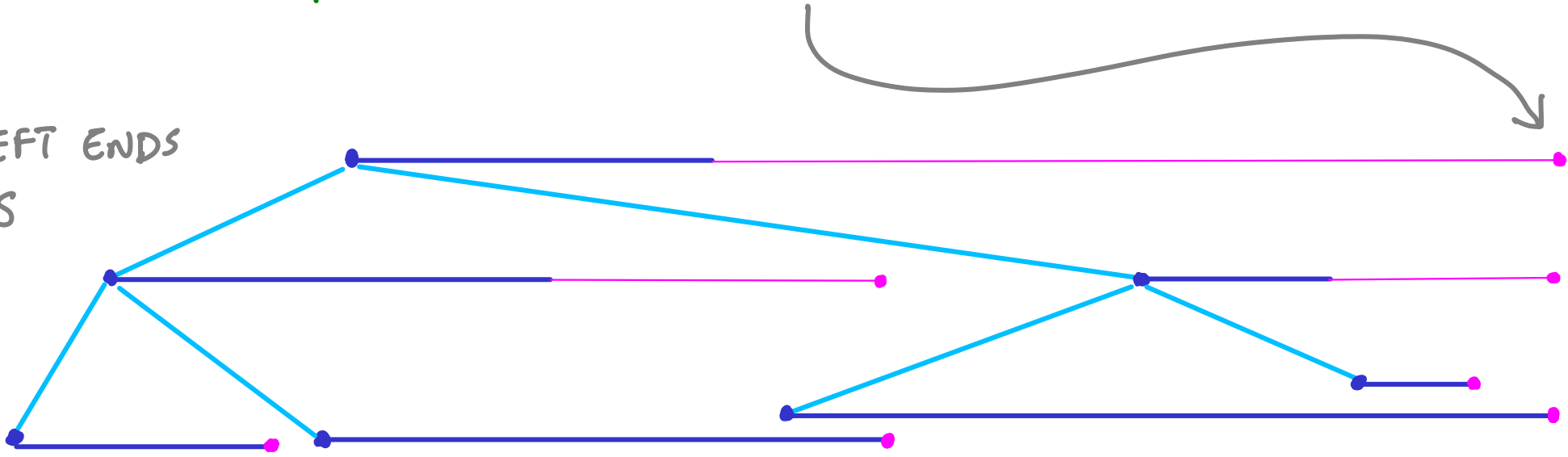
Spend  $O(1)$  time at a node,  
the recurse on only one side.

Given balanced tree, total time is  $O(\log n)$

But can we build and maintain the augmented info?

How can we update MAX RIGHT END OF SUBTREE ?

BST w/ LEFT ENDS  
as KEYS





7 \_\_\_\_\_ 10

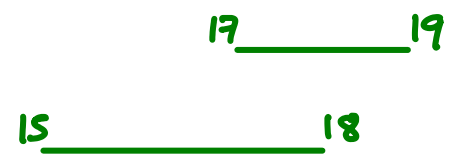
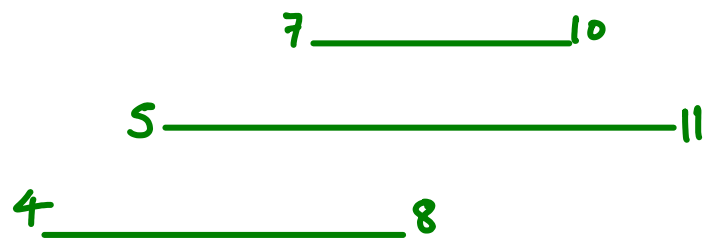
17 \_\_\_\_\_ 19

5 \_\_\_\_\_ 11

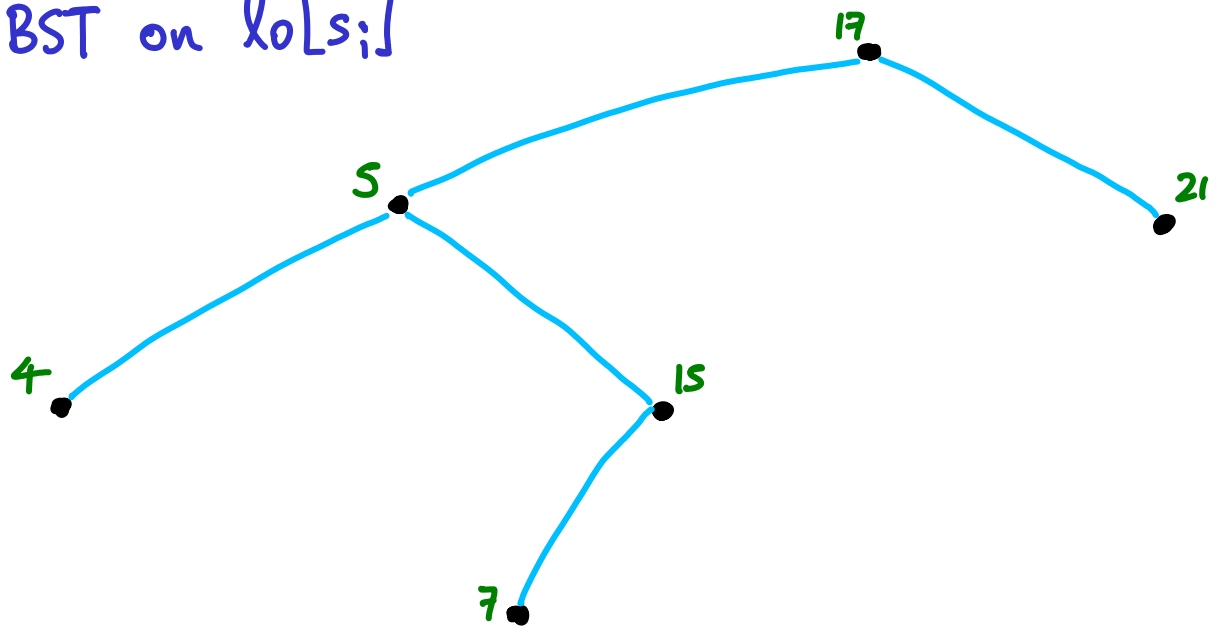
15 \_\_\_\_\_ 18

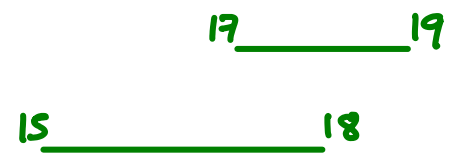
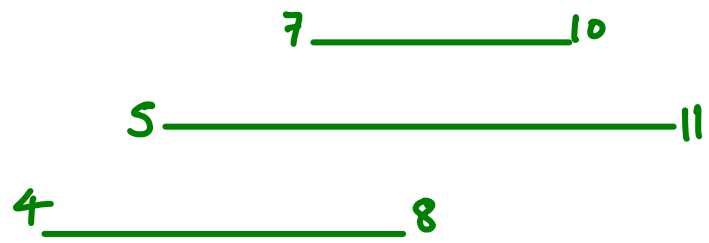
4 \_\_\_\_\_ 8

21 \_\_\_\_\_ 23

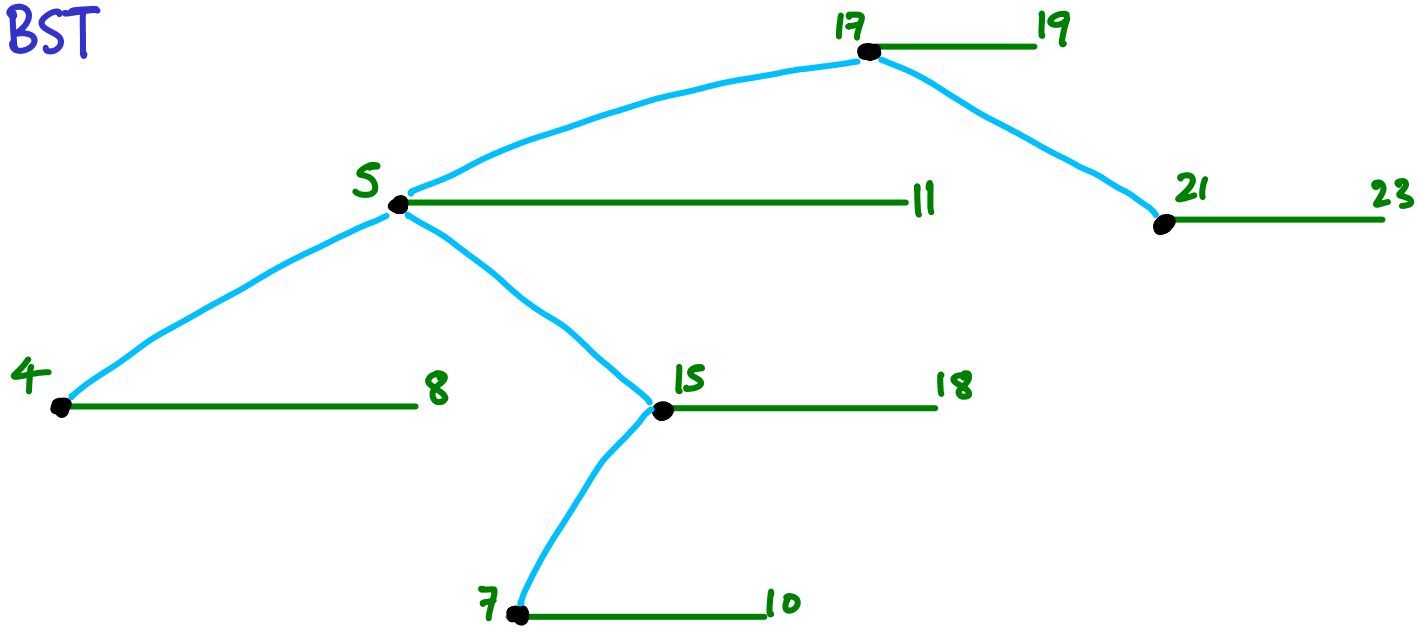


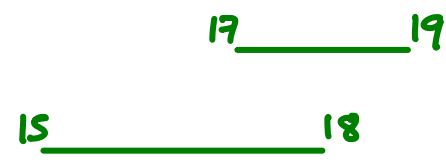
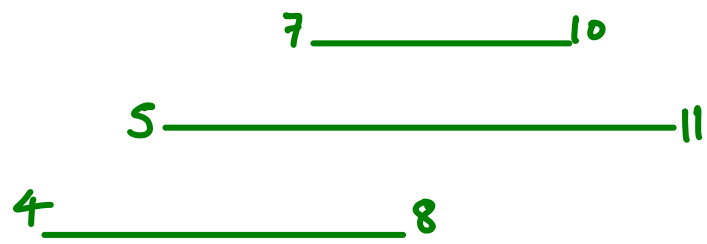
BST on  $lo[s;i]$



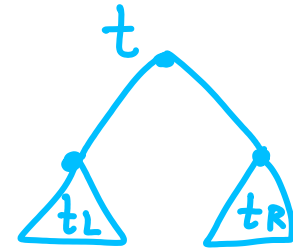
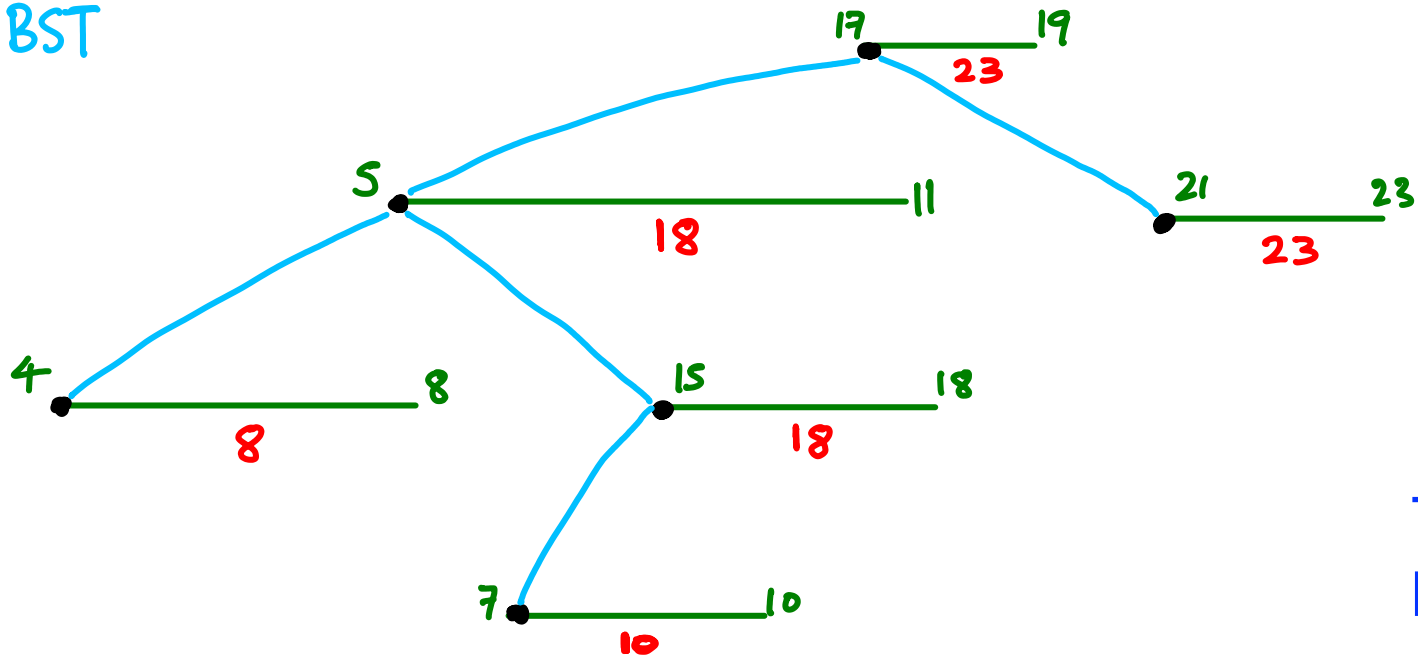


BST



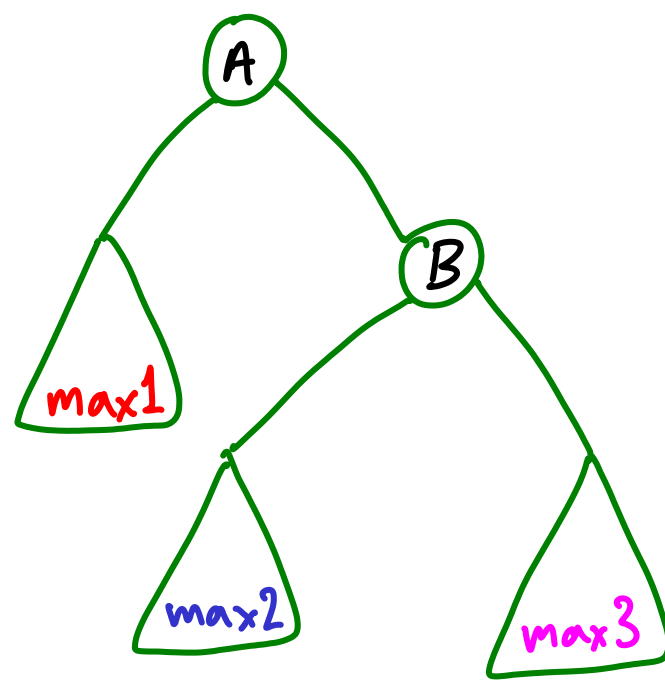
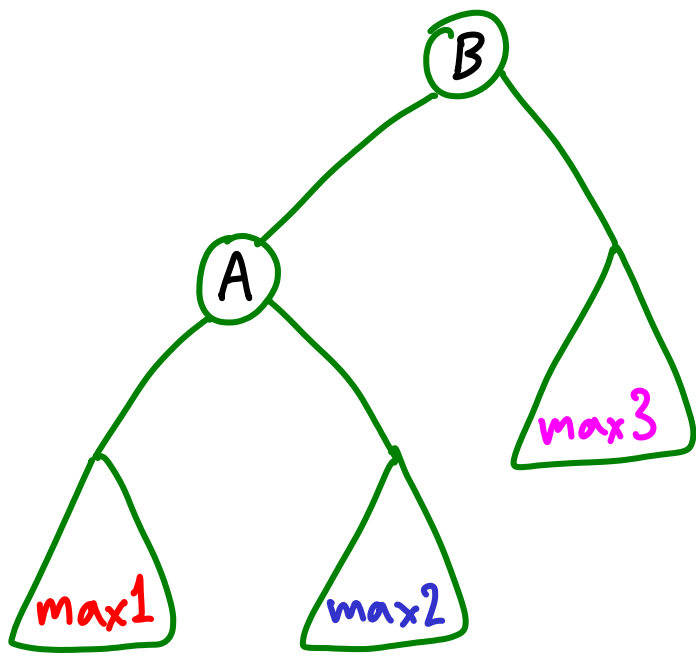


augmented  
BST



$$\max(t) = \max \begin{cases} h_i(t) \\ \max(t_L) \\ \max(t_R) \end{cases}$$

Trivial to build:  
bottom up after tree is made,  
or during insertion



$max1$ ,  $max2$ ,  $max3$  : unchanged by rotation  
 $max(A)$  &  $max(B)$  : trivial to update

we can maintain a balanced BST augmented w/ max value of subtrees