


DYNAMIC PROGRAMMING - LONGEST INCREASING SUBSEQUENCE

S: 23, 3, 5, 18, 10, 101, 12, 14, 4, 105

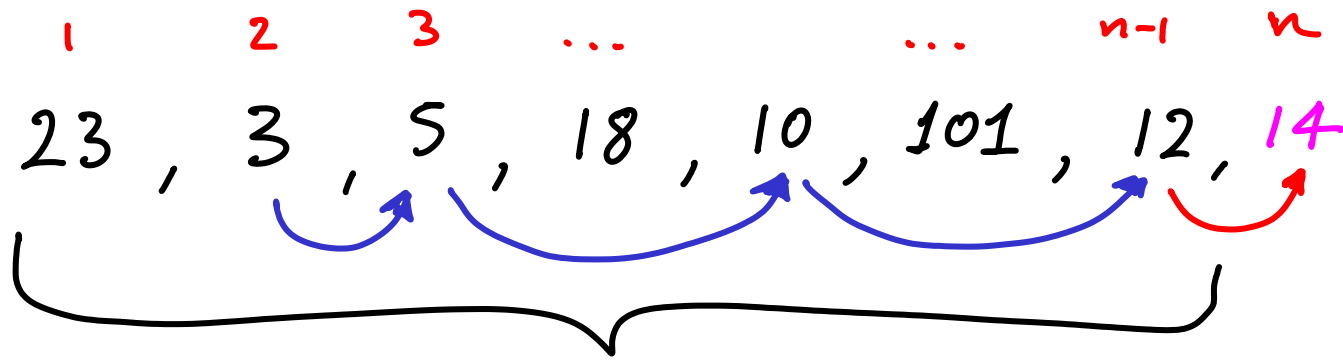


$L(S) = 3, 5, 10, 12, 14, 105$ $|L(S)| = 6$

Brute-force: try all subsequences, see if they are increasing: $O(2^n)$

For dynamic programming we would like

- a recursive expression w/ repeated subproblems
- an easy, fast way to use solved subproblems



$$L(S) \sim$$

$$\sim L_{1..n}(S)$$

$$\sim L_n$$

$|L_n|$ using $|L_{n-1}|$?


if $S[n] > \underline{\text{last element in } L_{n-1}}$ then $|L_n| = |L_{n-1}| + 1$
 ↳ could be at any position in S

else $\parallel S[n] \leq \text{last}(L_{n-1})$

↳ keep $|L_{n-1}|$?

↳ add $S[n]$ to suboptimal solution from $S[1..n-1]$?

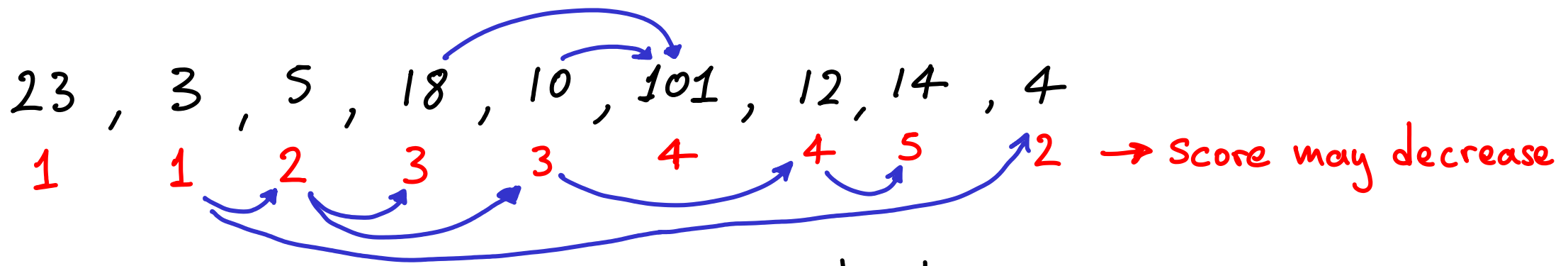


1 2 3 \dots \dots $n-1$ n
 23 , 3 , 5 , 18 , 10 , 101 , 12 , 14 , 4 , 105


Redefine: L_n = longest increasing subsequence
 that actually uses $S[n]$

$$|L_n| = 1 + \text{MAX}_{\{ \text{all } j \text{ s.t. } S[j] < S[n] \}} |L_j|$$

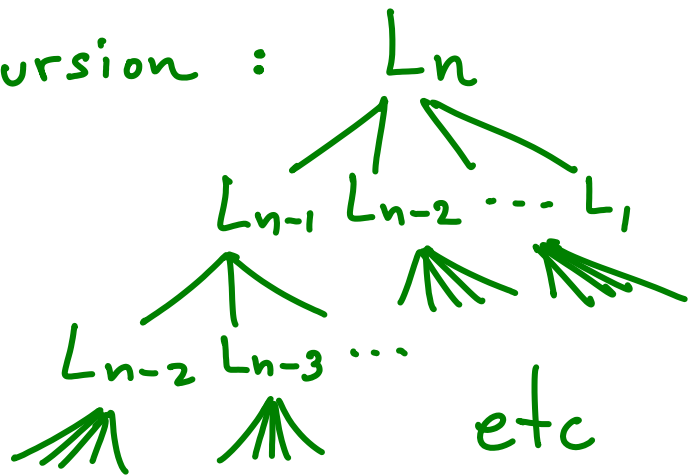
look at all L_j ($j < n$)



$$|L_n| = 1 + \text{MAX}_{\{ \text{all } j \text{ s.t. } S[j] < S[n] \}} |L_j|$$

Recursion :

BAD



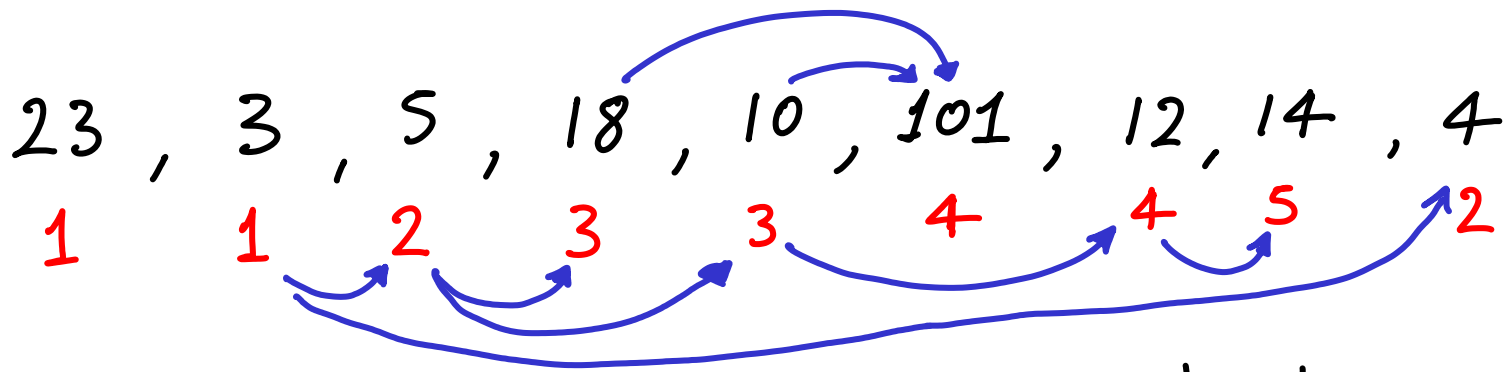
Dyn. Prog: Build solutions, "bottom up"

When it's time to solve $|L_k|$ we have stored all $|L_j|$ ($j < k$) in an array.

$$T(k) = \Theta(k)$$

$$T(n) = \sum_{i=1}^n T(k) = \Theta(n^2)$$

$$\text{Space} = \Theta(n)$$



$$T(n) = \Theta(n^2)$$

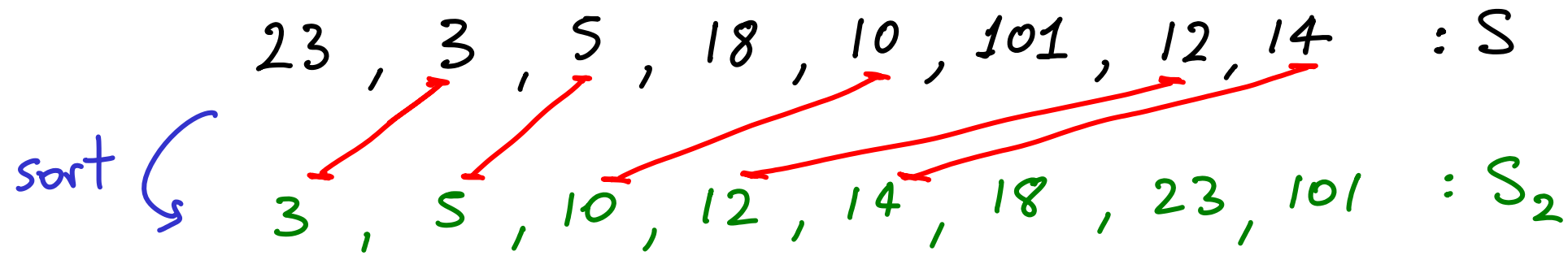
$$\text{space} = \Theta(n)$$

$$|L_n| = 1 + \text{MAX}_{\{ \text{all } j \text{ s.t. } S[j] < S[n] \}} |L_j|$$

What about $|L.I.S.|$? $= \max_{j=1..n} |L_j|$

What about L.I.S.? Keep the pointers: for each $S[j]$ store any $S[i]$ pointer that generated $|L_j|$

A QUICK SOLUTION FOR L.I.S. ... but still $O(n^2)$ & dyn-prog.



FIND LONGEST COMMON SUBSEQUENCE !

- any common subsequence is increasing
so $LCS(S, S_2)$ qualifies as a solution
- LIS must exist in S_2 , so it is a candidate for LCS.

□