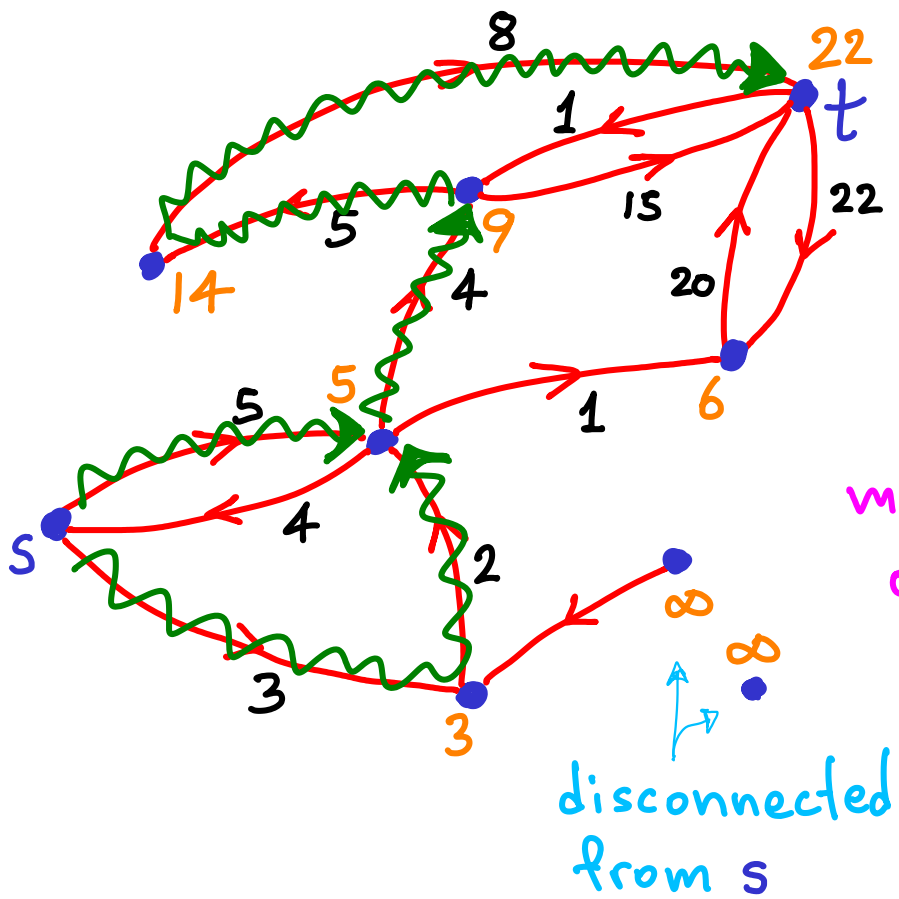


# SINGLE SOURCE SHORTEST PATHS



paths from  $s$  to  $t$

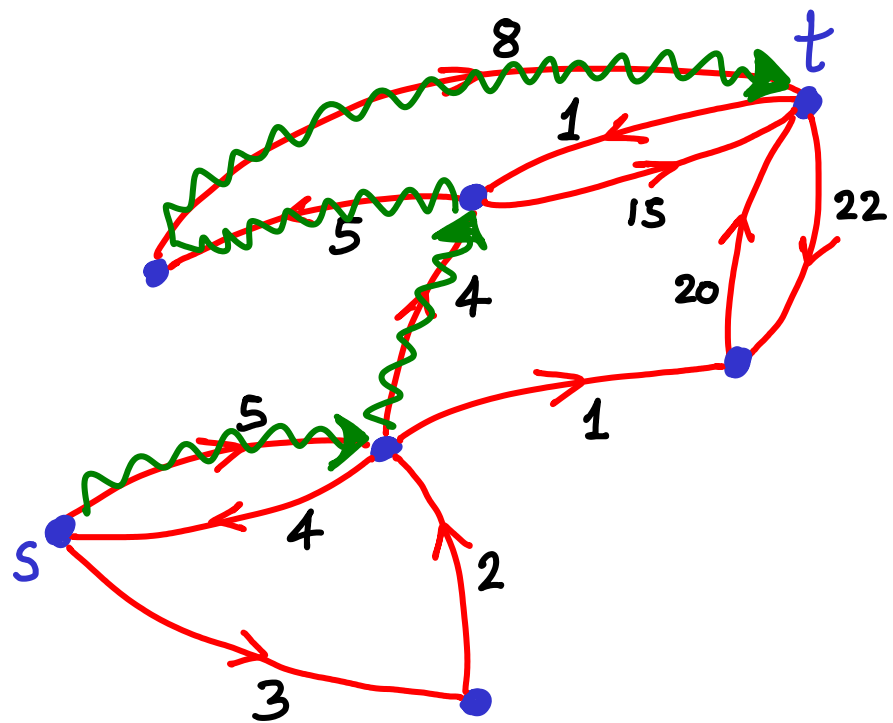
$$\begin{array}{r}
 3 + 2 + 1 + 20 \\
 3 + 2 + 4 + 15 \\
 \underline{3 + 2 + 4 + 5 + 8} = 22 \\
 \underline{5} + 1 + \dots \\
 \quad + 4 + \dots = 22
 \end{array}$$

not greedy  
or BFS

multiple options {

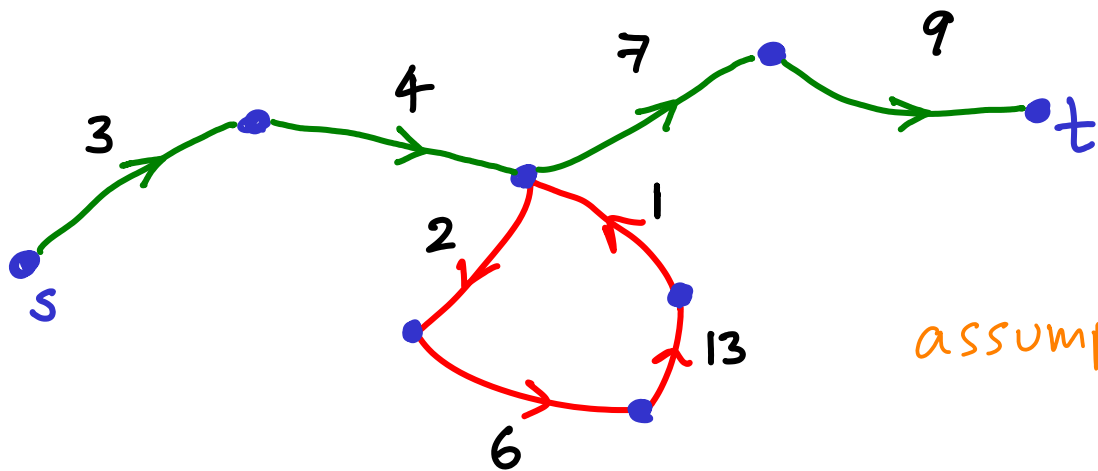
Generally assume a directed graph  
(can make undirected  $\rightarrow$  directed easily)

# SINGLE SOURCE SHORTEST PATHS



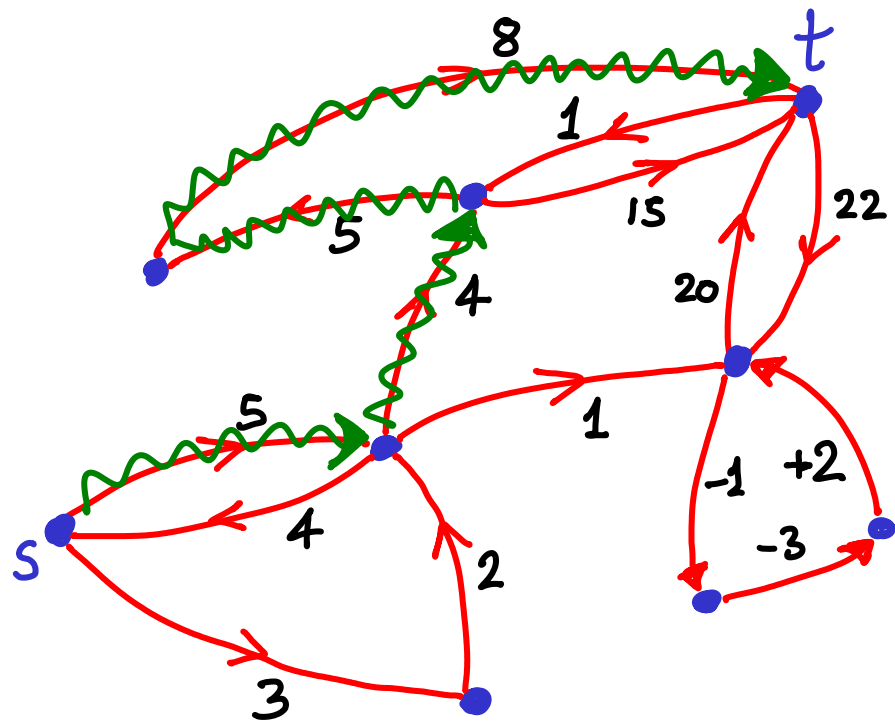
## Observations

- No cycles in  $s \rightarrow t$  (shortest path)



assumption?

# SINGLE SOURCE SHORTEST PATHS



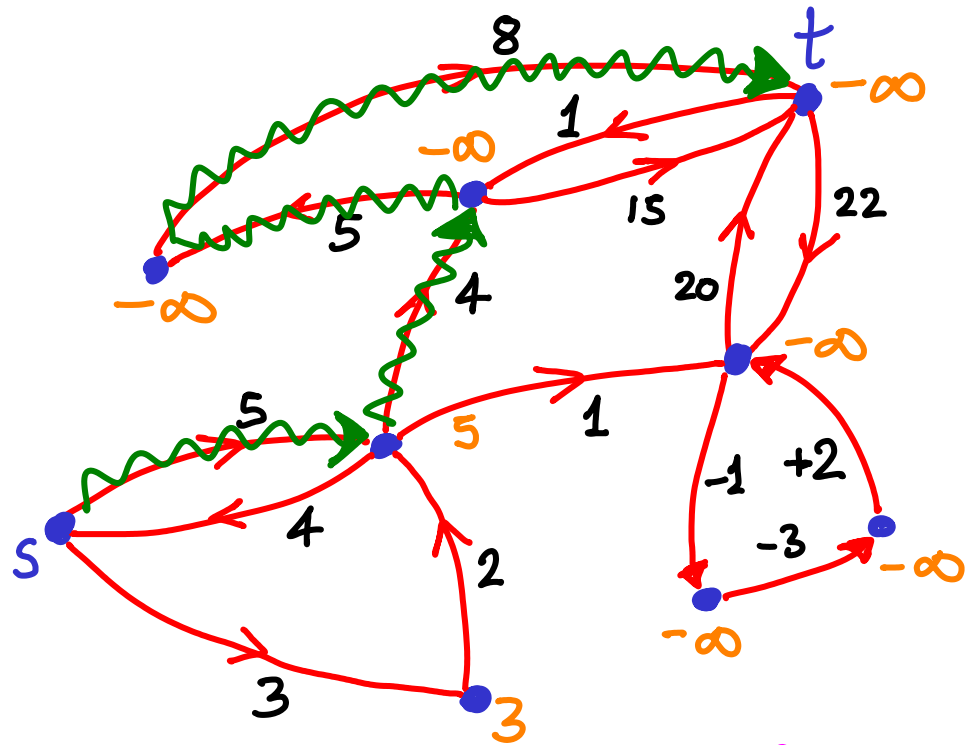
## Observations

- No cycles in  $s \rightarrow t$
- Negative weights ~OK, unless they form a **negative cycle** in  $G$

$$\Sigma(\text{cycle}) < 0$$

# SINGLE SOURCE

# SHORTEST PATHS



## Observations

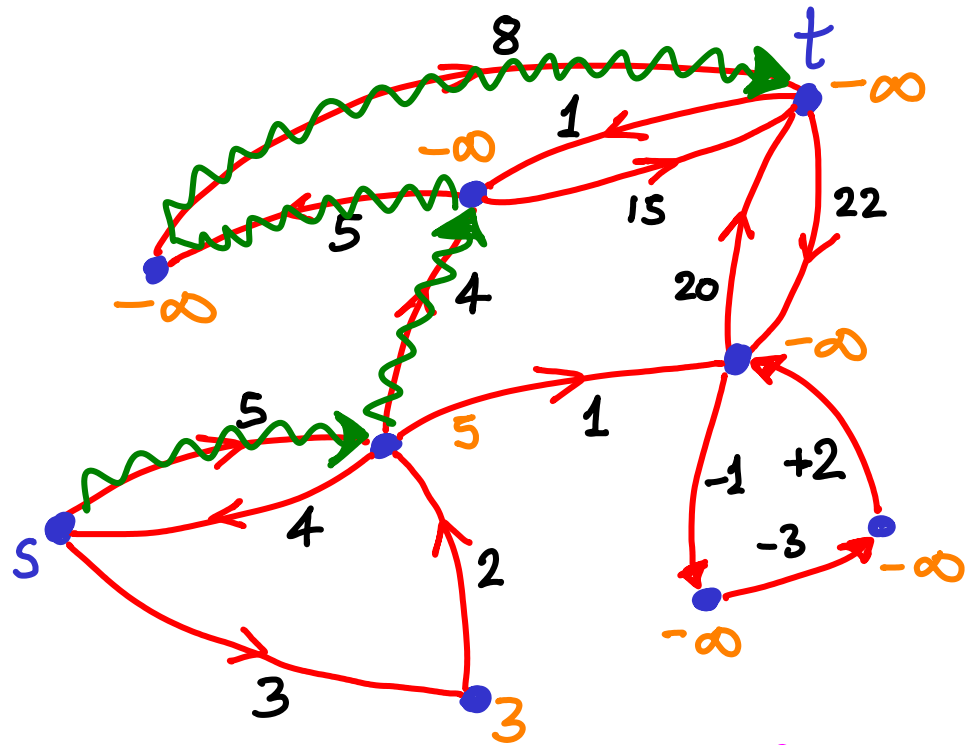
- No cycles in  $s \rightarrow t$
- Negative weights ~OK, unless they form a **negative cycle** in  $G$

Any vertex reachable from a negative cycle gets a score of  $-\infty$

} assuming cycle can be reached from  $s$

# SINGLE SOURCE

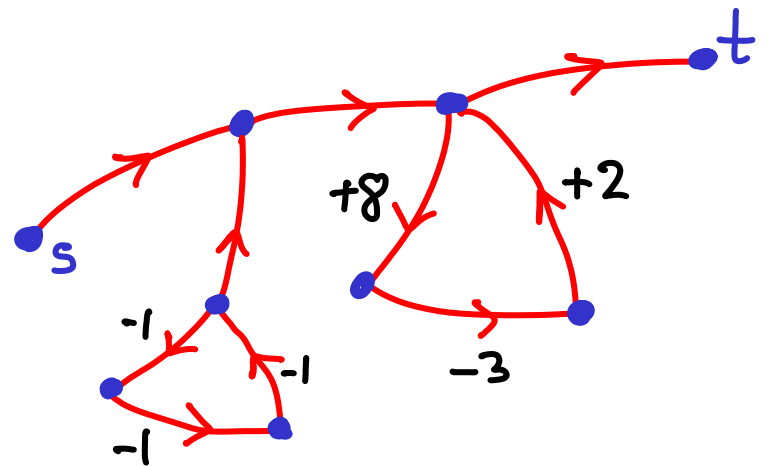
# SHORTEST PATHS



Any vertex reachable from a negative cycle gets a score of  $-\infty$

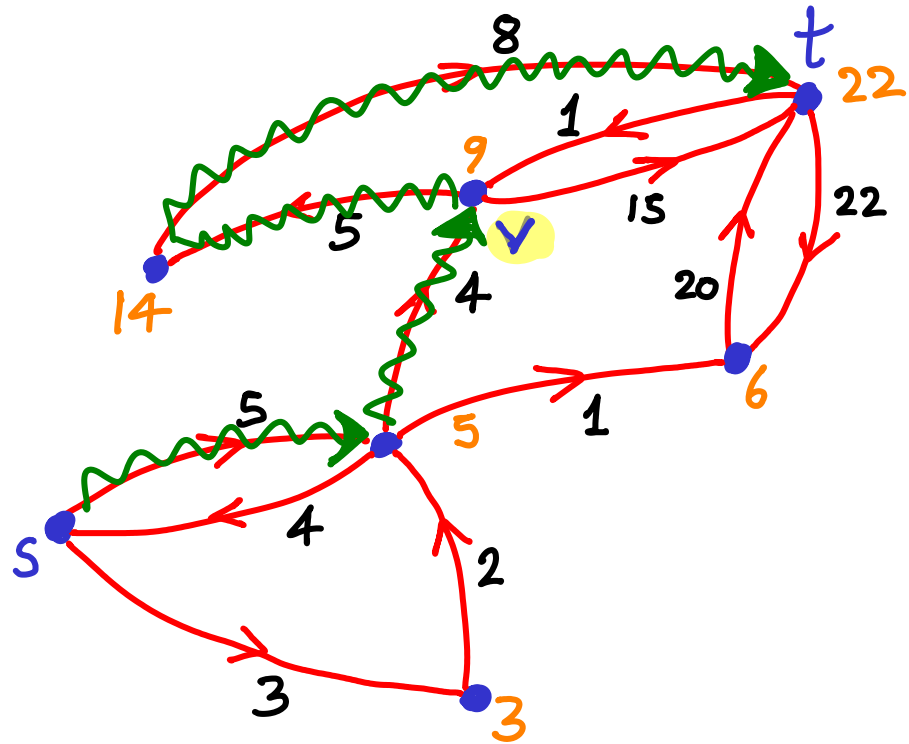
## Observations

- No cycles in  $s \rightarrow t$
- Negative weights ~OK, unless they form a **negative cycle** in  $G$



# SINGLE SOURCE

# SHORTEST PATHS



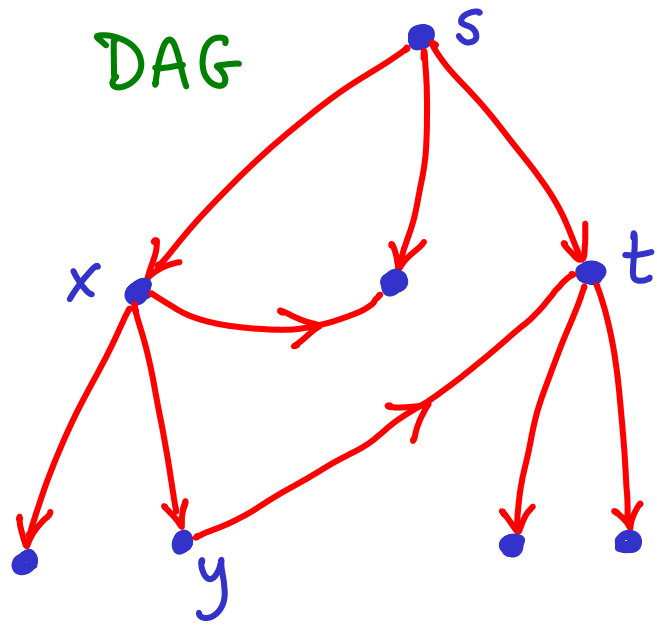
## Observations

- No cycles in  $s \rightarrow t$
- Negative weights ~OK, unless they form a **negative cycle** in  $G$
- shortest path  $s \rightarrow v \rightarrow t$  contains

shortest path  $s \rightarrow v$  (9)

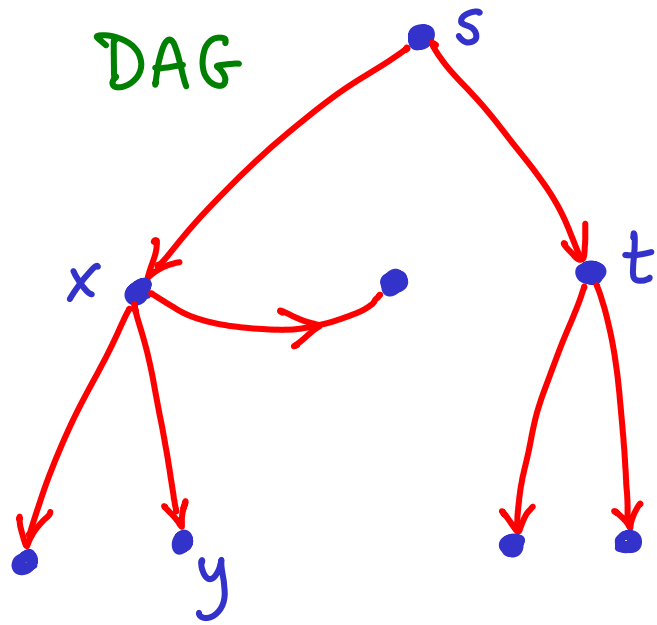
&

shortest path  $v \rightarrow t$  (13)



there may be multiple shortest paths  
e.g.  $s \rightarrow t$  or  $s \rightarrow x \rightarrow y \rightarrow t$

All shortest paths from  $s$  to  $V$   
can be represented in a DAG



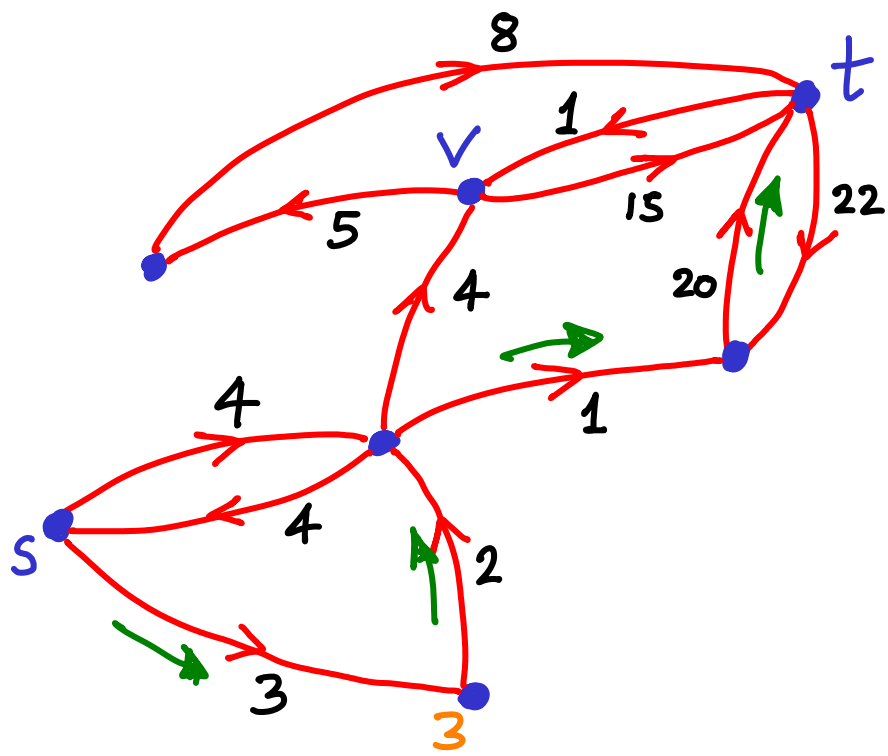
there may be multiple shortest paths  
e.g.  $s \rightarrow t$  or  $s \rightarrow x \rightarrow y \rightarrow t$

All shortest paths from  $s$  to  $V$   
can be represented in a DAG

DAG  $\rightarrow$  tree : arbitrarily keep one path to each vertex  
"shortest paths tree"

(similar to picking one BFS/DFS search)





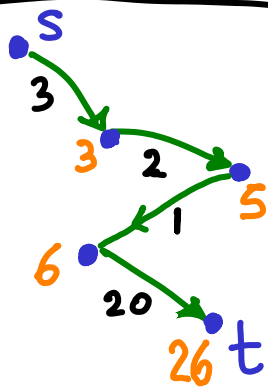
By exploring some path from  $s$  to  $t$  we get a path score (e.g. 26)

the score of  $t$  is 26, which may only decrease as we explore more options.

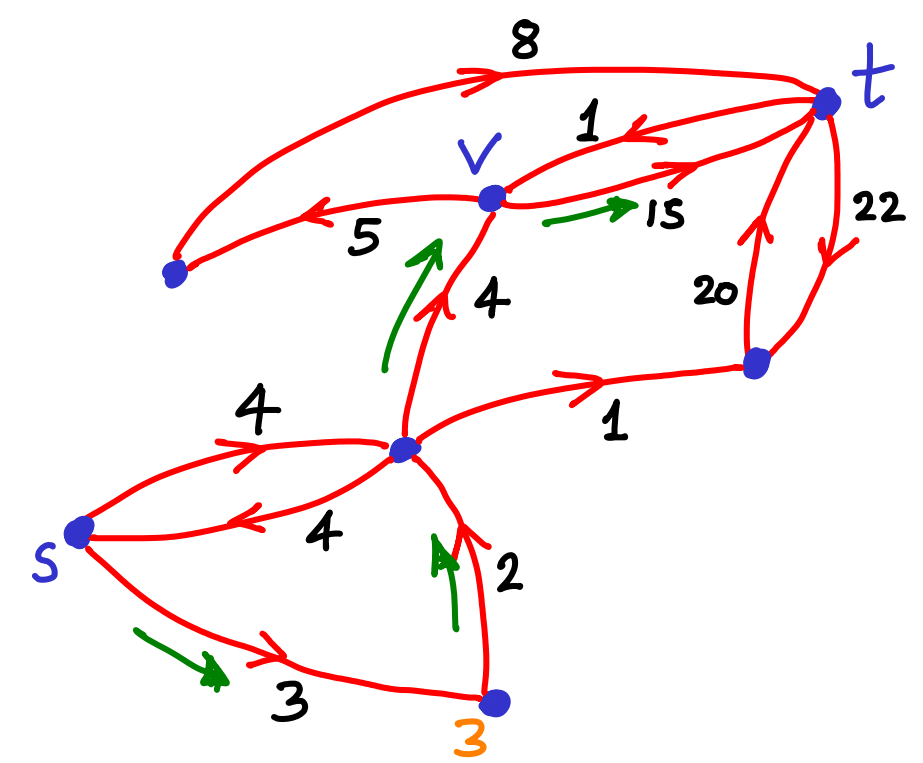
If we update the score of  $v$ :  $d(v)$  &  $\exists$  edge  $v \rightarrow t$

then we can possibly improve  $d(t)$ :  
 $d(v) + w(v, t) < d(t)$ ?  
 (15)

shortest paths tree



change tree



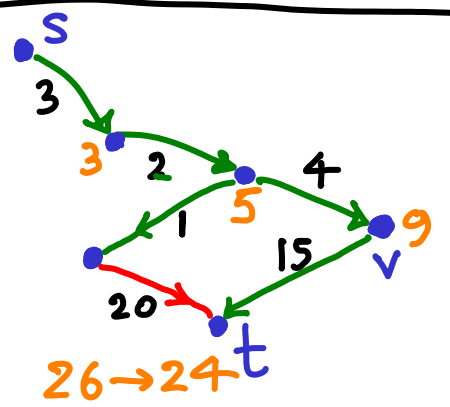
By exploring some path from  $s$  to  $t$  we get a path score (e.g. 26)

the score of  $t$  is 26, which may only decrease as we explore more options.

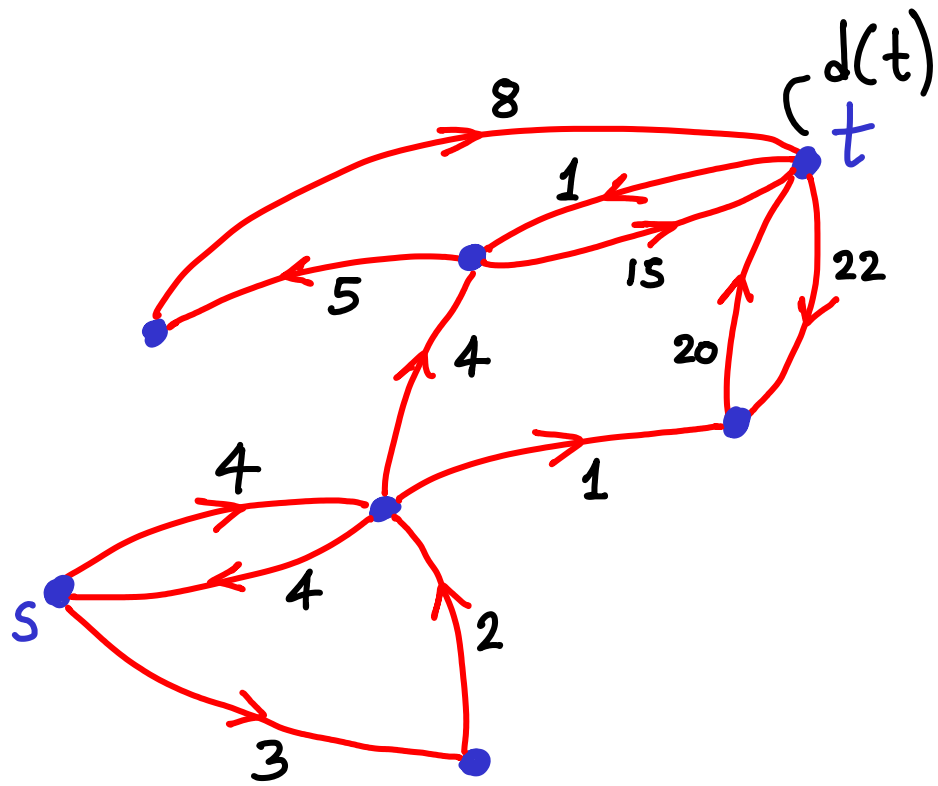
If we update the score of  $v$ :  $d(v)$  &  $\exists$  edge  $v \rightarrow t$

then we can possibly improve  $d(t)$ :  
 $d(v) + w(v, t) < d(t)$ ?  
 (15)

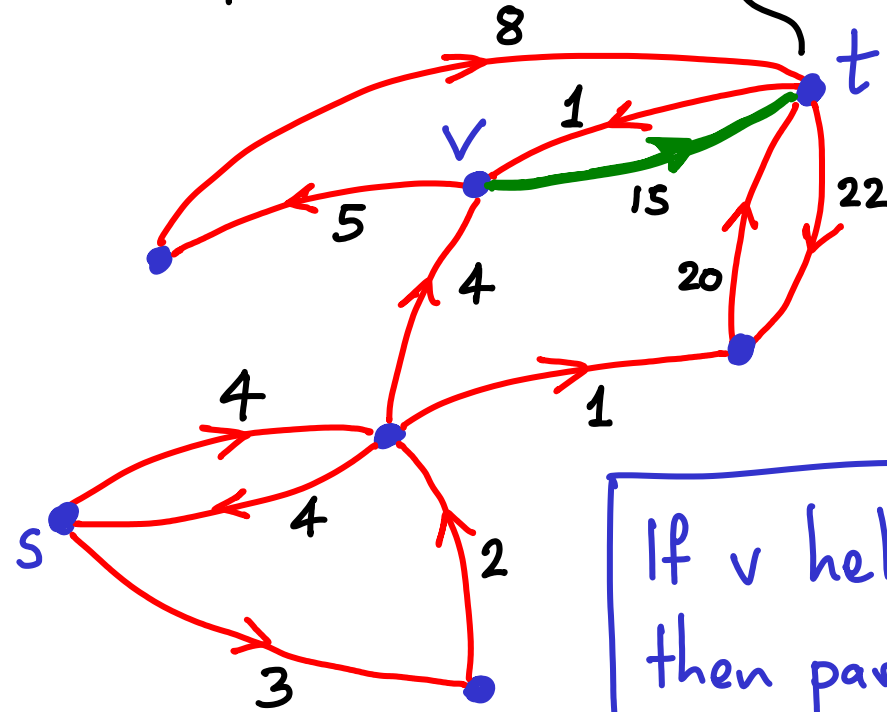
shortest paths tree



change tree



Keep MIN of  $d(t)$  vs.  $d(v) + w(v, t)$

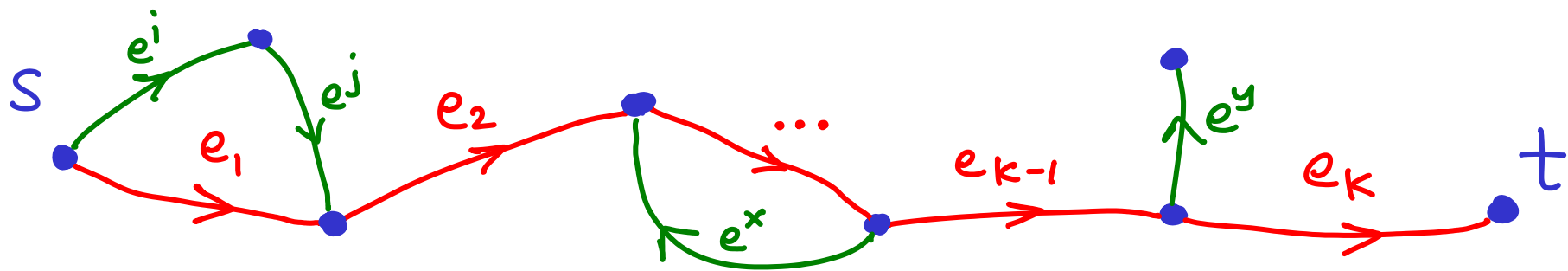


If  $v$  helps,  
then  $\text{parent}(t) = v$

**Relax( $v, t$ )**

: checking if score of  $t$  can be improved (lowered) by using  $s \rightarrow v \rightarrow t$

Assume this is a shortest path from  $s$  to  $t$  unknown but exists



Suppose we have an algorithm based on relaxing edges.

If we relax  $e_1$  before  $e_2$  before  $\dots$  before  $e_{k-1}$  before  $e_k$

then we will correctly compute  $d(t)$  by INDUCTION

Relax sequence :  $e^x e_1 e^j e^y e_2 e^x e^i e_k e_{k-1} e_1 e^x e_k e^y$  : OK  
 (don't care if we relax other edges or the same ones repeatedly)