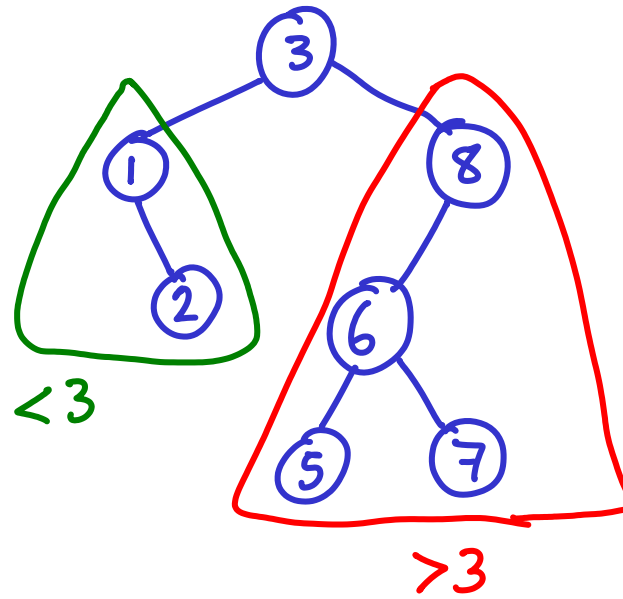
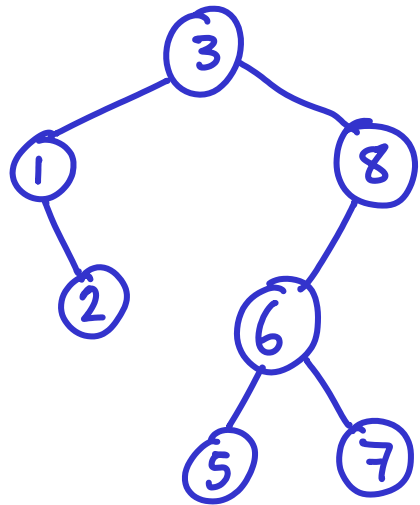
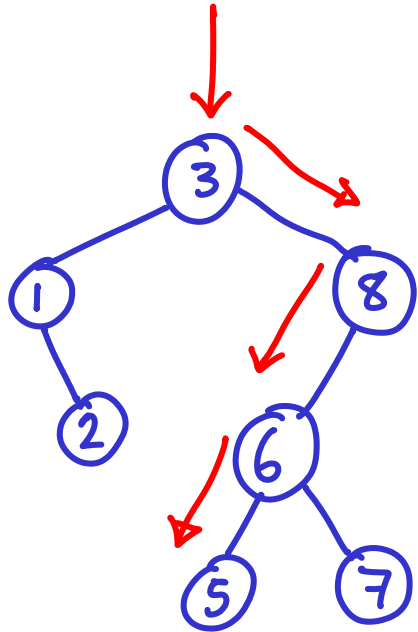


# BINARY SEARCH TREES



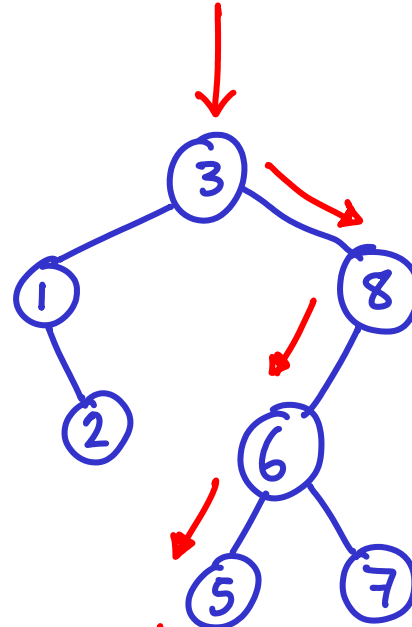
# BINARY SEARCH TREES

(binary) search for 5



found

(binary) search for 4



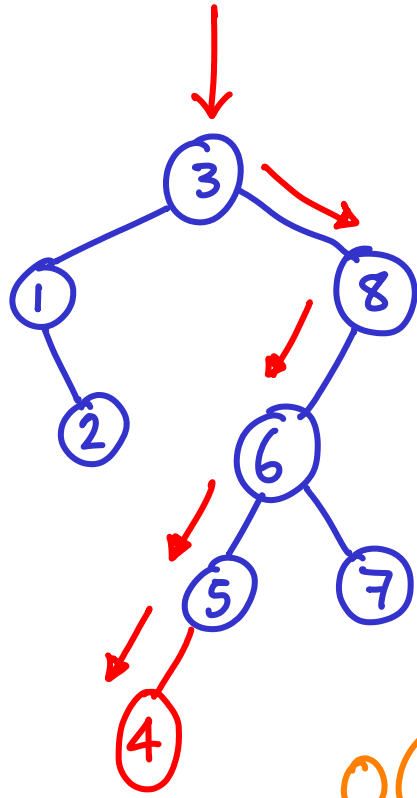
not found



time:  
 $O(\text{depth})$

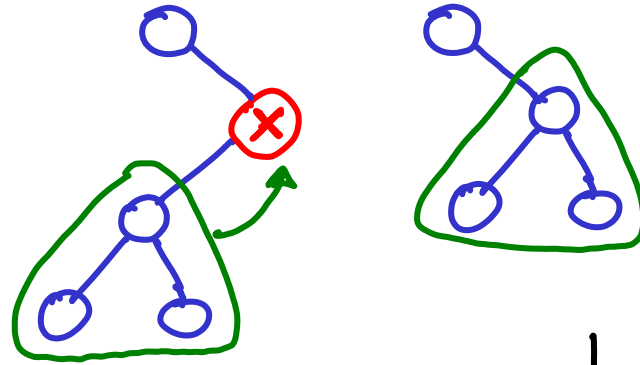
# BINARY SEARCH TREES

insert (4) ~ search



$O(\text{depth})$

instant delete (x)



x: any node w/  $< 2$  children

If (one) subtree exists,  
promote it.

$O(1)$

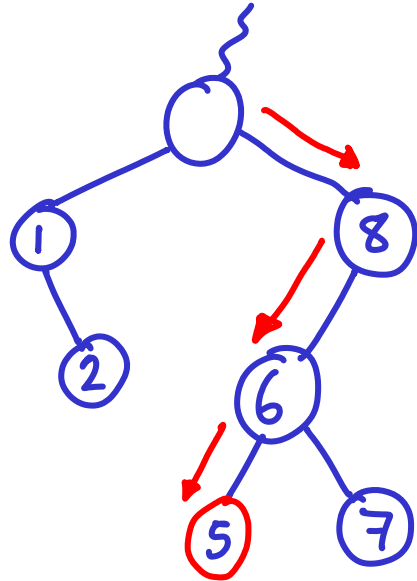
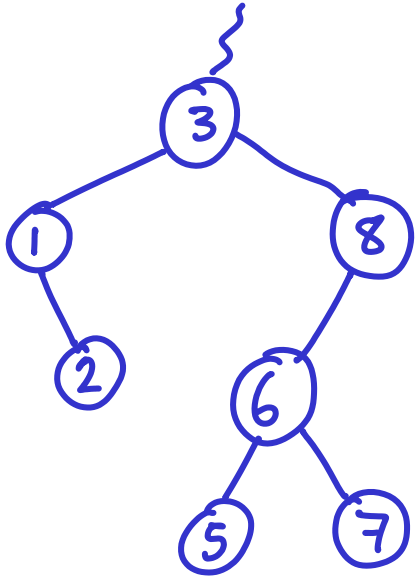
# BINARY SEARCH TREES

non-instant  
delete(3)



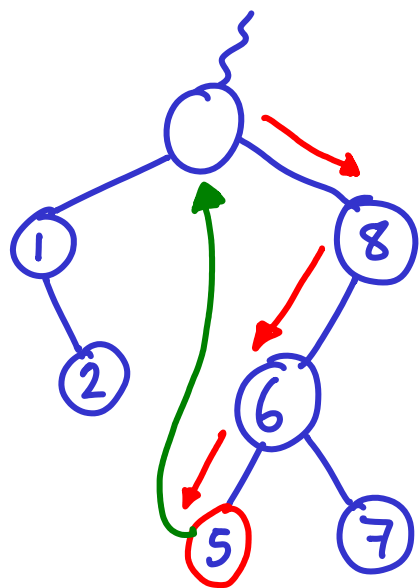
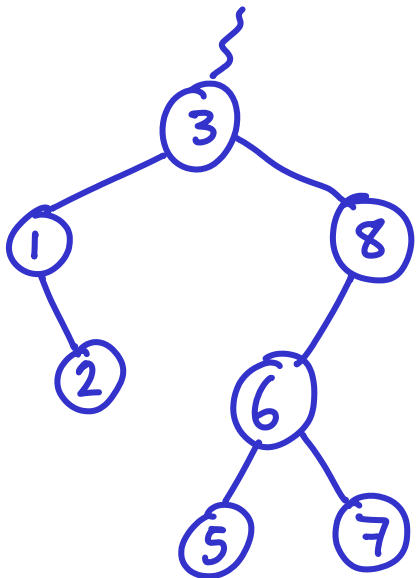
find successor

: smallest element greater than 3  
(which exists because : 2 children)



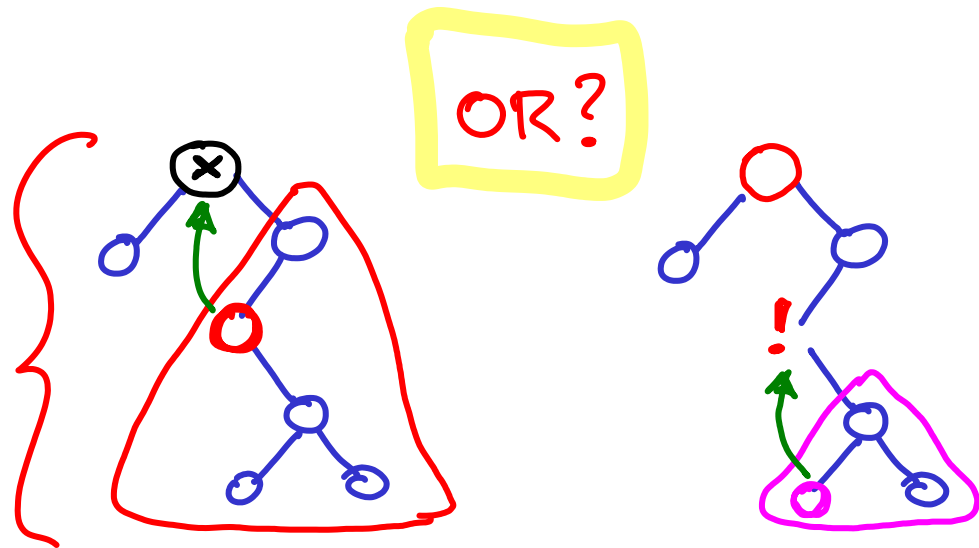
# BINARY SEARCH TREES

delete(3) → find successor & replace



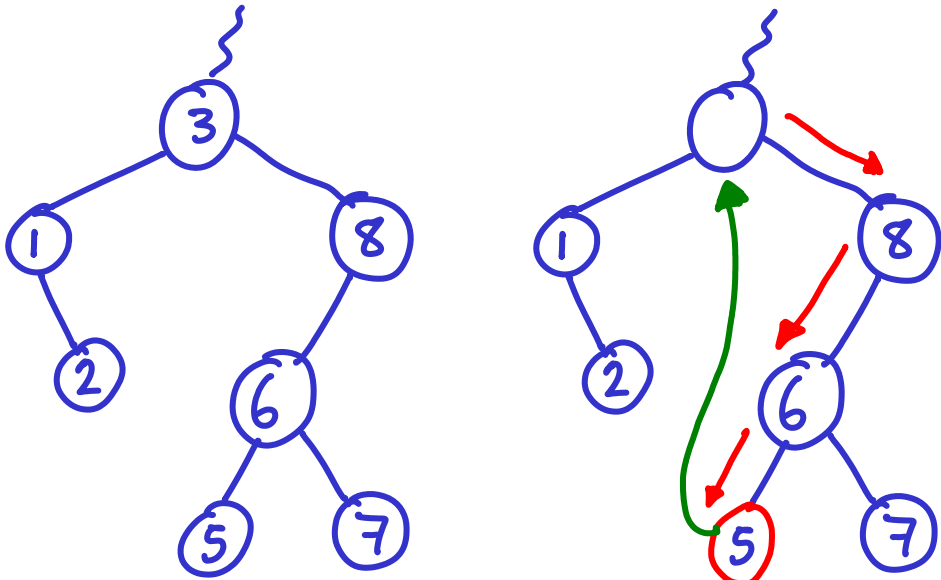
total  $O(\text{depth})$  ← might need to recurse

By definition, successor is the last node visited on a path from R-child(3)



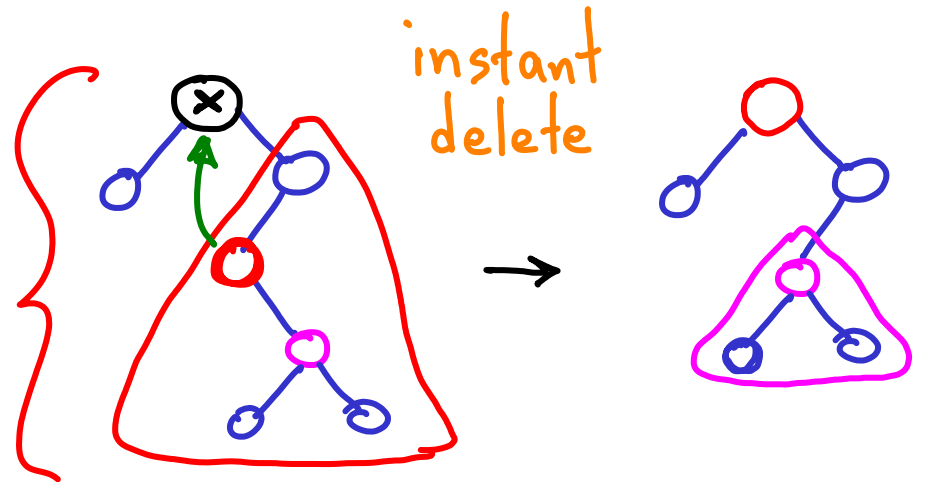
# BINARY SEARCH TREES

delete(3) → find successor & replace

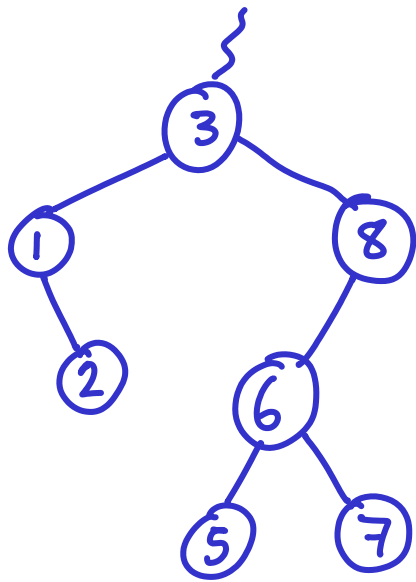


still  $O(\text{depth})$  ← no need to recurse

By definition, successor is the last node visited on a ↙ path from R-child(3)



# BINARY SEARCH TREE SUMMARY



INSERT

DELETE

SEARCH →  $O(\text{depth})$

We should keep the tree balanced  
as much as possible